

# **An Algebra for Assessing Trust in Authentication Chains**

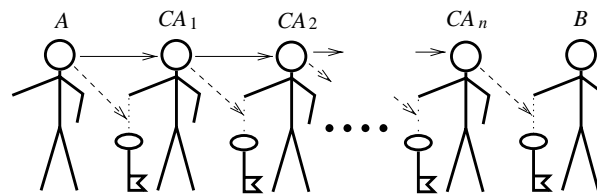
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# Key authenticity based on chains of trust

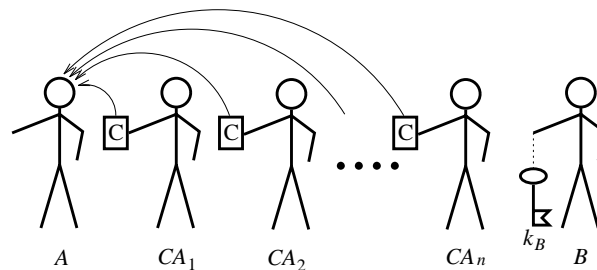
Two types of trust:

1. Trust in key authenticity (key-to-owner binding).
2. Recommendation trust (agent co-operation).



Legend:  
——> Recommendation Trust (RT)  
-----> Trust in Key Authenticity (KA)

Agent *A* must determine the authenticity of  $k_B$  based on recommendations in the form of certificates.



## The belief model

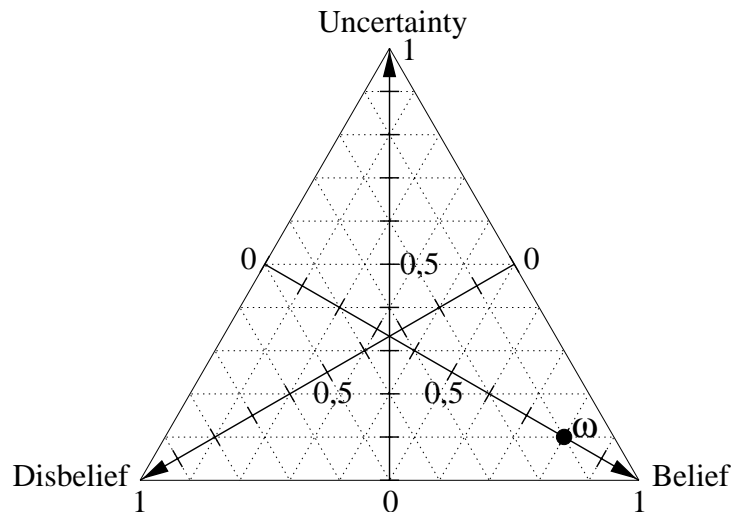
An *opinion* is a triplet  $\{b, d, u\}$  which satisfies

$$b + d + u = 1, \quad \{b, d, u\} \in [0, 1]^3$$

$b$ : belief

$d$ : disbelief

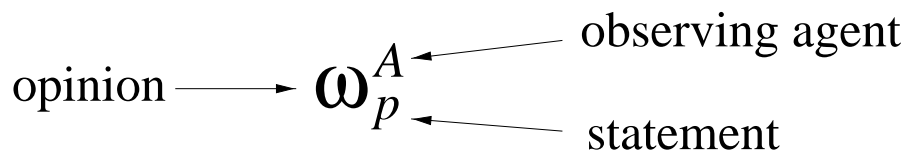
$u$ : uncertainty



- Any point in the triangle represents an opinion.
- Example,  $\omega = \{0.8, 0.1, 0.1\}$  is represented as a dot in the triangle.

## Subjective Logic

- Opinions can be interpreted as imprecise probabilities of binary events.
- Subjective Logic is reduced to probability calculus when  $u = 0$ .
- Subjective logic is reduced to binary logic when  $b = 1$  or  $d = 1$ .
- Ownership of opinions is assigned to individuals.



## The operators of Subjective Logic

1. AND  $\omega_p^A \wedge \omega_q^A$

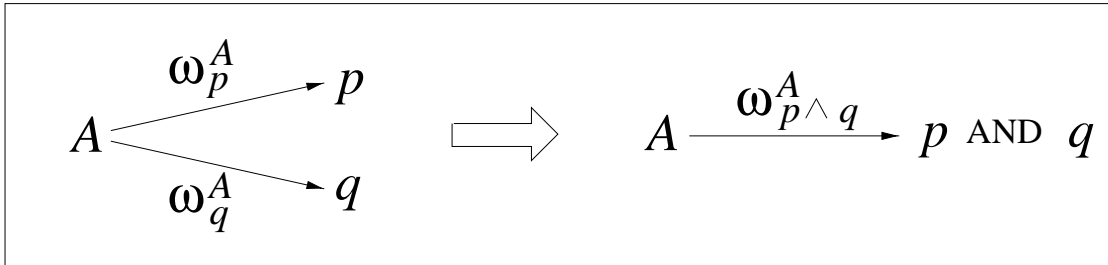
2. OR  $\omega_p^A \vee \omega_q^A$

3. Negation  $\neg\omega_p^A$

4. Recommendation  $\omega_p^A \otimes \omega_p^B$

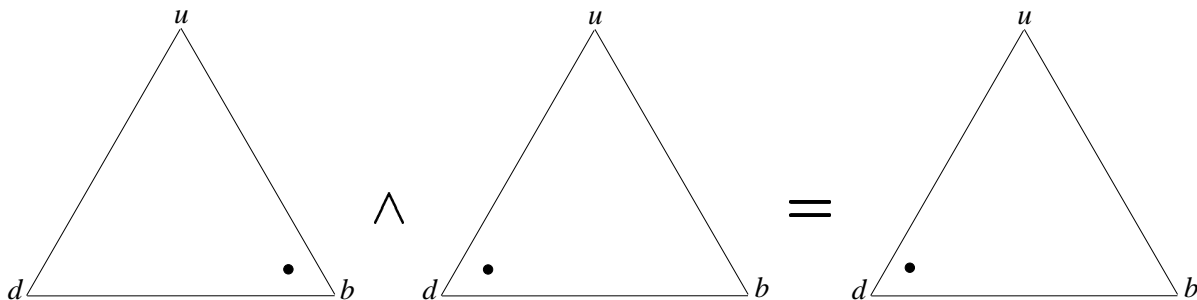
5. Consensus  $\omega_p^A \oplus \omega_p^B$

## AND

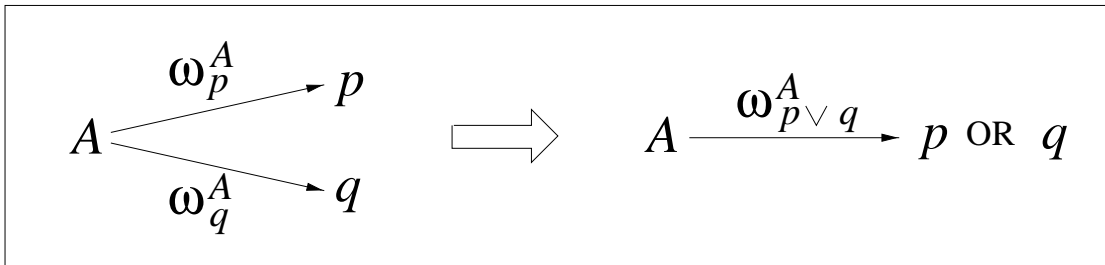


- notation:  $\omega_p^A \wedge \omega_q^A = \omega_{p \wedge q}^A$
- commutative
- associative
- opinion independence assumed
- not idempotent:  $\omega_p^A \wedge \omega_p^A$  is undefined
- becomes product of probabilities i.c.o. zero ignorance
- becomes 'binary logic AND' i.c.o. absolute opinions

Ex:  $\{0.8, 0.1, 0.1\} \wedge \{0.1, 0.8, 0.1\} = \{0.08, 0.82, 0.10\}$

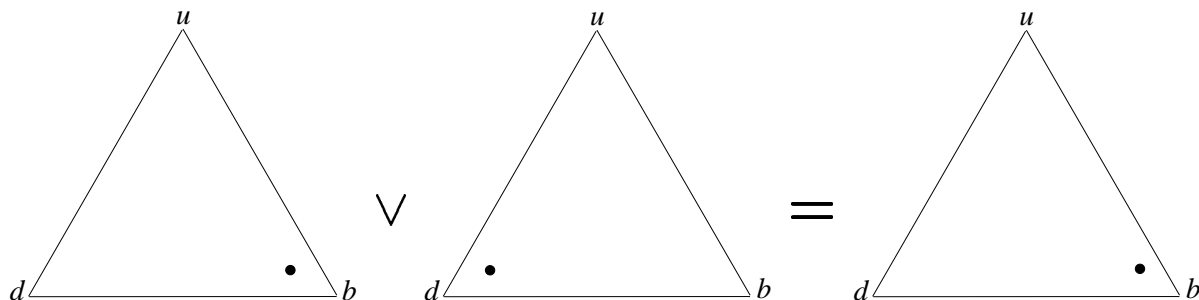


# OR

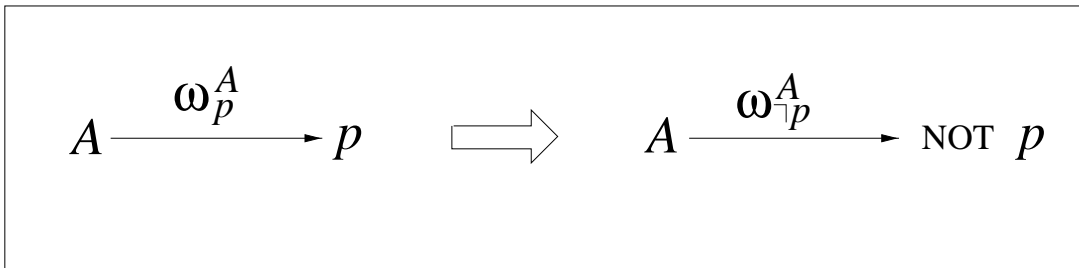


- notation:  $\omega_p^A \vee \omega_q^A = \omega_{p \vee q}^A$
- commutative
- associative
- opinion independence assumed
- not idempotent:  $\omega_p \vee \omega_p$  is undefined
- becomes co-product of probabilities i.c.o. zero ignorance
- becomes 'binary logic OR' i.c.o. absolute opinions

Ex:  $\{0.8, 0.1, 0.1\} \vee \{0.1, 0.8, 0.1\} = \{0.82, 0.08, 0.10\}$

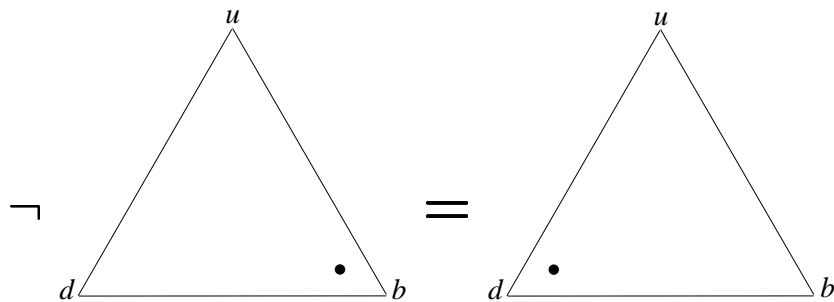


## Negation



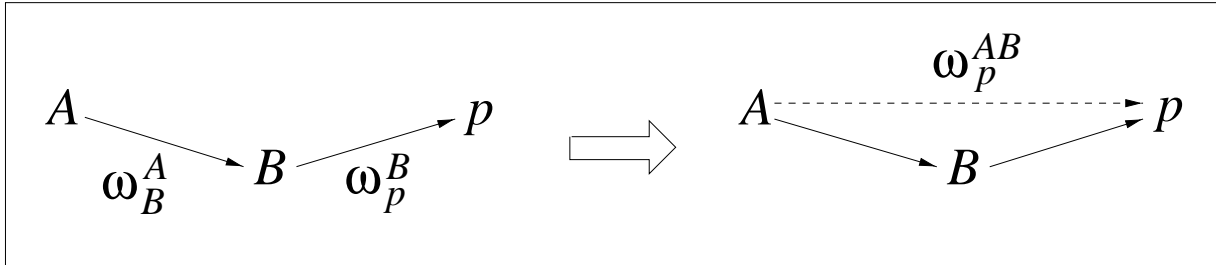
- notation:  $\neg \omega_p^A = \omega_{\neg p}^A$
- Negation is involutive so that  $\neg(\neg \omega_p^A) = \omega_p^A$

Ex:  $\neg\{0.8, 0.1, 0.1\} = \{0.1, 0.8, 0.1\}$



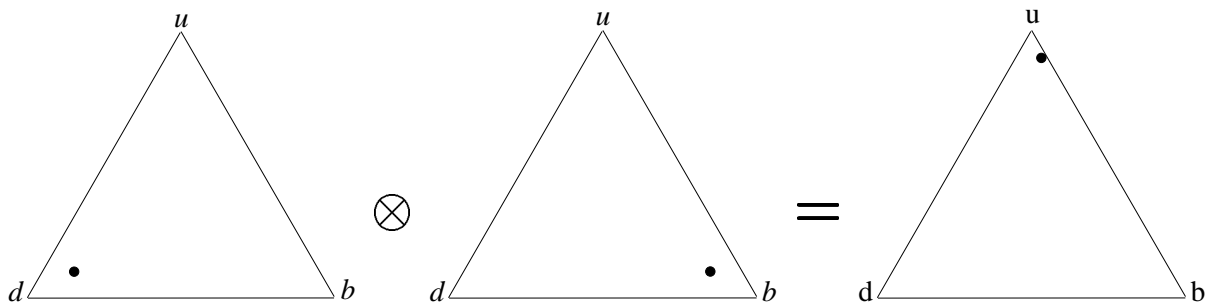


## Recommendation

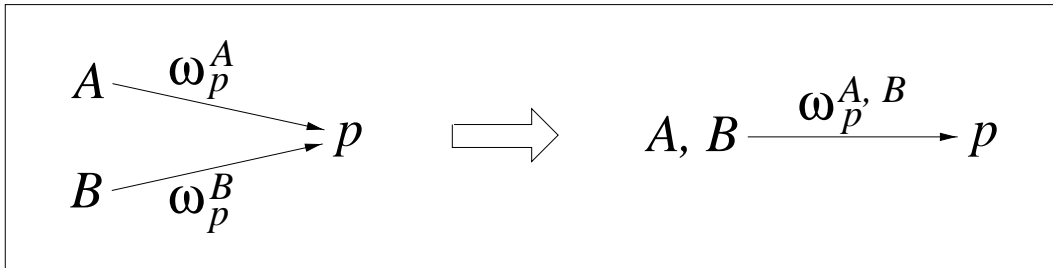


- notation:  $\omega_p^{AB} = \omega_B^A \otimes \omega_p^B$
- associative
- non-commutative
- opinion independence assumed
- transitivity assumed

Ex:  $\{0.1, 0.8, 0.1\} \otimes \{0.8, 0.1, 0.1\} = \{0.08, 0.01, 0.91\}$

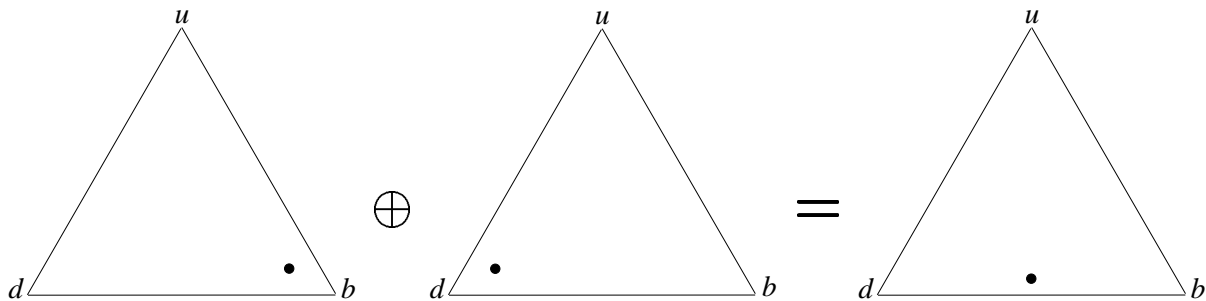


## Consensus



- notation:  $\omega_p^{A,B} = \omega_p^A \oplus \omega_p^B$
- commutative
- associative
- opinion independence assumed
- opinions without ignorance can not be combined

Ex:  $\{0.8, 0.1, 0.1\} \oplus \{0.1, 0.8, 0.1\} = \{0.47, 0.47, 0.06\}$

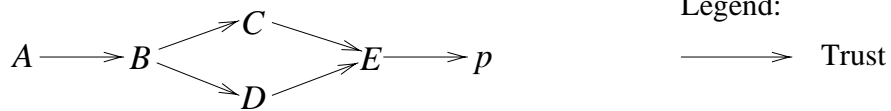


## The problem of dependence

‘AND’ and ‘OR’ are not distributive on each other:

$$\omega_p \wedge (\omega_q \vee \omega_r) \neq (\omega_p \wedge \omega_q) \vee (\omega_p \wedge \omega_r)$$

Recommendation is not distributive on consensus:



$$\begin{aligned} & \omega_B^A \otimes ((\omega_C^B \otimes \omega_E^C) \oplus (\omega_D^B \otimes \omega_E^D)) \otimes \omega_p^E \\ & \neq \\ & (\omega_B^A \otimes \omega_C^B \otimes \omega_E^C \otimes \omega_p^E) \oplus (\omega_B^A \otimes \omega_D^B \otimes \omega_E^D \otimes \omega_p^E) \end{aligned}$$

## Modelling trust

$p$ : “*The system will resist malicious attacks.*”

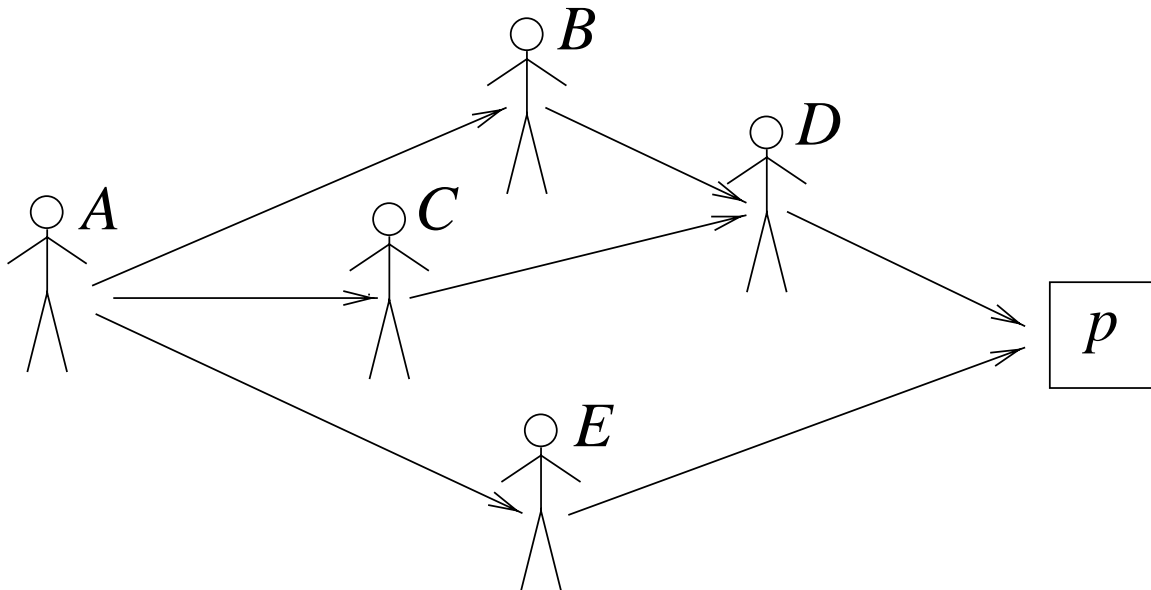
$q$ : “*The agent will cooperate.*”

$r$ : “*The key is authentic.*”

$\omega_p$ ,  $\omega_q$ , and  $\omega_r$  are trust parameters.

Trust models can be constructed using subjective logic.

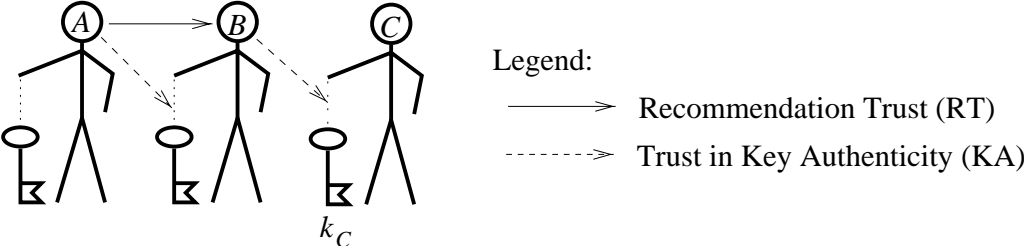
## Propagation of trust in social networks



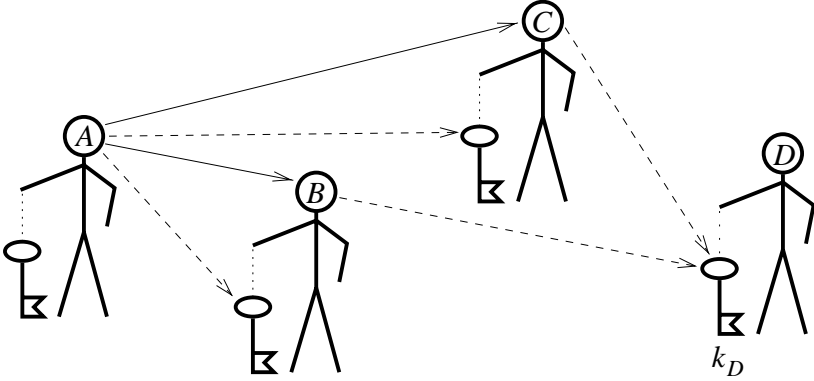
$$\omega_p^{(AB, AC)D, AE} = ((\omega_B^A \otimes \omega_D^B) \oplus (\omega_C^A \otimes \omega_D^C) \otimes \omega_p^D) \oplus (\omega_E^A \otimes \omega_p^E)$$

# Computation of key authenticity based on trust

Notation:  $\omega_B^A = (\omega_{RT(B)}^A \wedge \omega_{KA(k_B)}^A)$

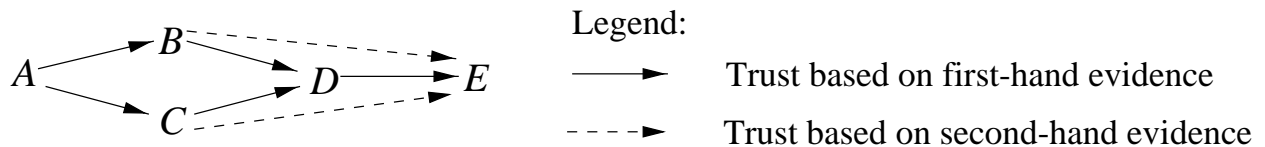


$$\omega_{k_C}^{AB} = \omega_B^A \otimes \omega_{k_C}^B$$



$$\omega_{k_D}^{AB,AC} = (\omega_B^A \otimes \omega_{k_D}^B) \oplus (\omega_C^A \otimes \omega_{k_D}^C)$$

## Warning: First-hand trust only!



If  $B$  and  $C$  recommends their second-hand trust to  $A$ , then  $A$  would think:

$$\omega_E^{AB,AC} = (\omega_B^A \otimes \omega_E^B) \oplus (\omega_C^A \otimes \omega_E^C)$$

Whereas in reality  $A$  would compute:

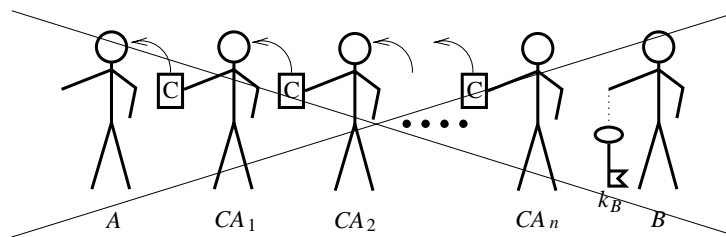
$$\omega_E^{ABD,ACD} = (\omega_B^A \otimes \omega_D^B \otimes \omega_E^D) \oplus (\omega_C^A \otimes \omega_D^C \otimes \omega_E^D)$$

The correct way is to recommend first-hand trust only:

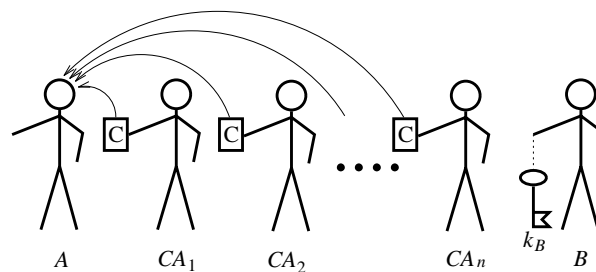
$$\omega_E^{(AB,AC)D} = ((\omega_B^A \otimes \omega_D^B) \oplus (\omega_C^A \otimes \omega_D^C)) \otimes \omega_E^D$$

## Direct routing of certificates

Indirect routing and re-computation of trust would lead to recommendation of second-hand trust.



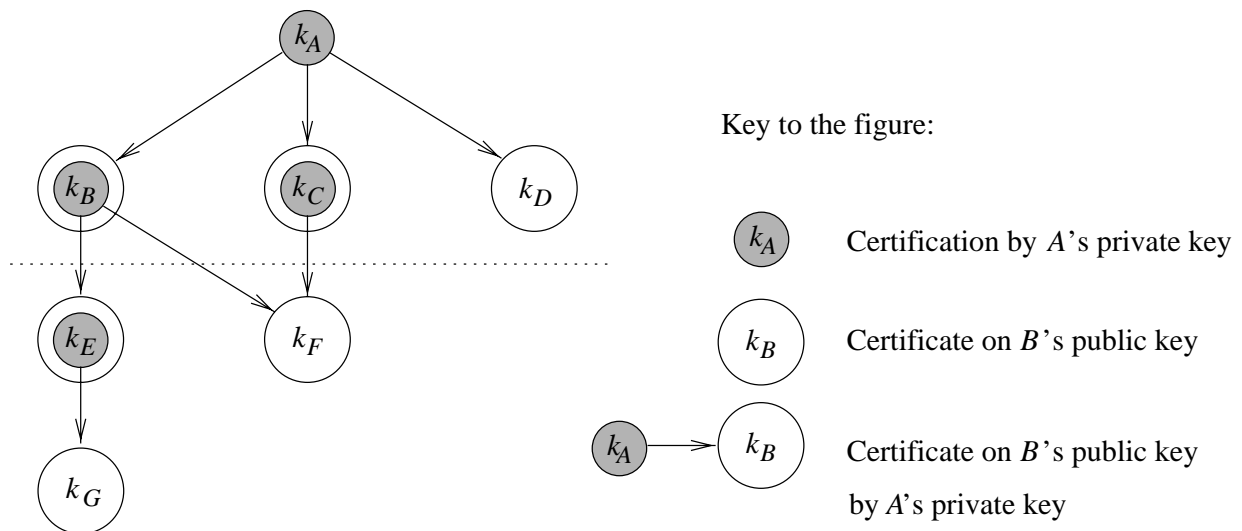
Recommendation of first-hand trust requires direct routing to the final recipient.



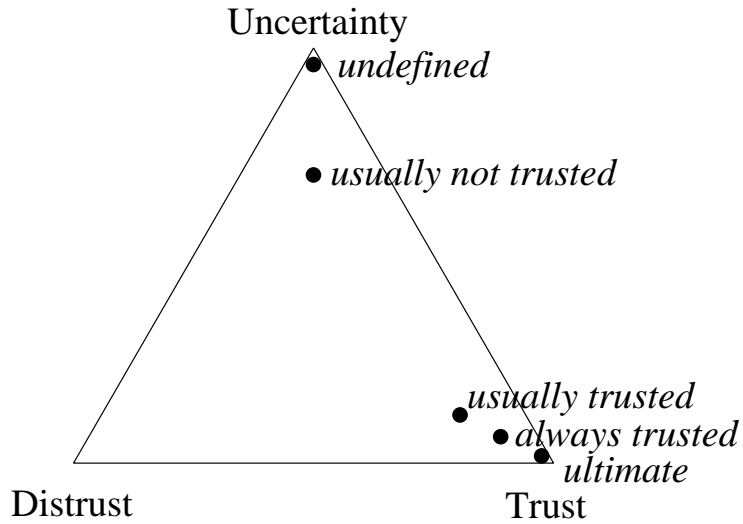


## Building a database of certified keys

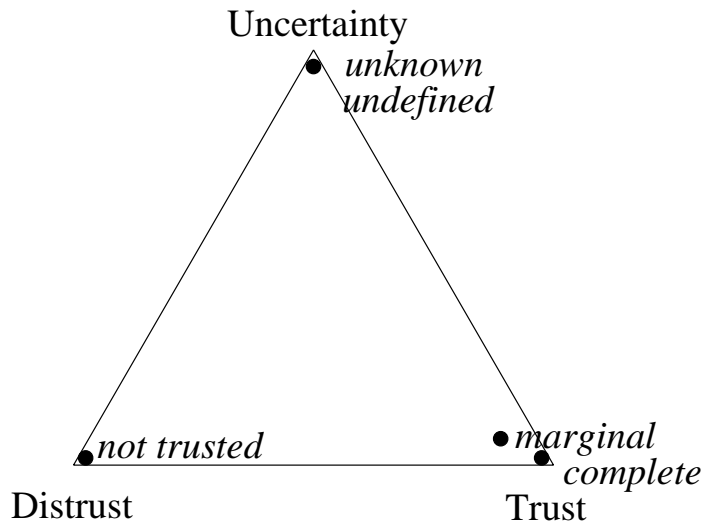
- Public keys can be exchanged manually or electronically.
- Electronically received keys must be certified.
- Each agent decides which other agents she will trust.



# Expressing PGP trust values

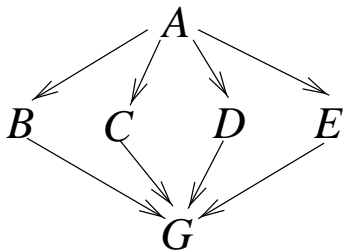


a) "Owner Trust" and "Signature Trust"

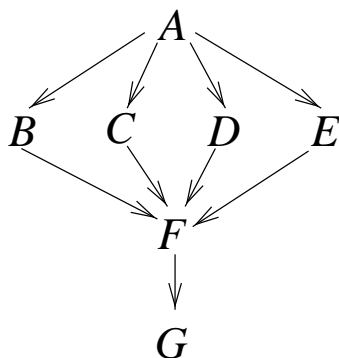


b) "Key Legitimacy"

## Hidden dependencies in PGP trust values



a) The situation that A sees



b) The real situation which is hidden for A

A thinks  $\omega_{KA(k_G)}^{AB,AC,AD,AE}$ ,

but computes  $\omega_{KA(k_G)}^{ABF,ACF,ADF,AEF}$ .

A should have computed  $\omega_{KA(k_G)}^{(AB,AC,AD,AE)F}$ .

## Concluding remarks

- The presented trust model is more complete than previously proposed models because it can express degrees of uncertainty.
- Subjective Logic can be used directly for reasoning about trust in practical security applications.
- A key certificate must contain:
  - 1) recommendation about key authenticity,
  - 2) recommendation about key owner.
- Recommendation of trust must be based on first-hand evidence only.