

Trust Models In ICE-TEL

Andrew Young
Nada Kapidzic Cicovic
David Chadwick

**Interworking Public Key
Certification Infrastructure for
Europe**



European 4th Framework Project under the TELEMATICS Programme





Overview

- A quick look at public key authentication
- Comparison of existing trust models
- ICE-TEL, the best of both worlds
- Examples

Public Key Authentication

- To verify a digital signature, I need
 - the signer’s public key
 - to be sure who “owns the public key”
 - » (i.e. who knows the corresponding private key)
- Certification
 - Third party assertion of “who owns which public key”
- Which third parties do I trust?
 - On what basis do they make their assertion?
 - What guarantees do they give? Liability?



Certification

- Third party issues certificate, comprising:
 - Who is doing the asserting (issuer)
 - Who is the subject of the assertion
 - What is being asserted (public key)
 - The small print (certification policy)
 - Digital signature
- Syntactic check of certificates tells me if the public key is accurate
- Semantic check of policies tells me who the public key belongs to
 - or what can be done with it

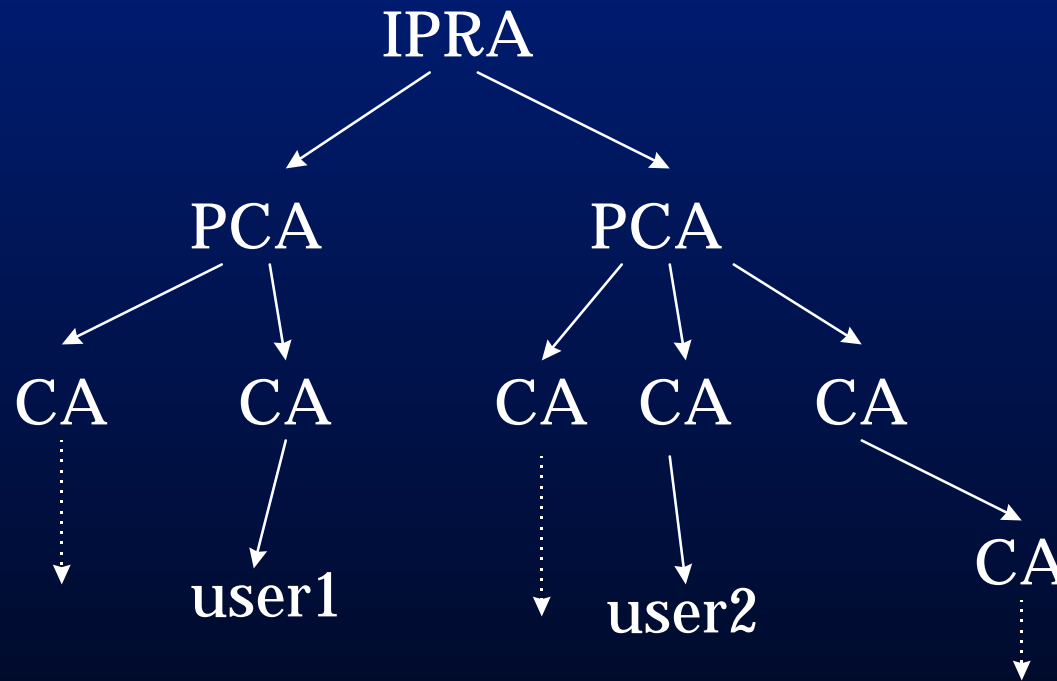
PGP Trust Model

- Web of trust
- Third party is “Trusted Introducer”
- Introducer does not have a “policy”



PEM Trust Model

- CA hierarchy
- PCAs publish a “certification policy”
- IPRA ties the PCAs together





The Gap in the Market

- PGP is user-centric
- PGP does not scale up to large communities
- PEM is organisation-centric
- PEM does not scale down to small communities



The ICE-TEL Trust Model

- Supports diverse security domains
 - single users
 - simple groups or small organisations
 - complex organisations
- Supports organic growth, allowing reorganisation of domains
- Trust between domains is by choice, and need not be mutual or transitive
- No central infrastructure



Trust Points

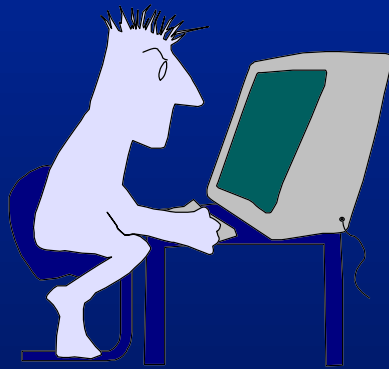
- Each security domain contains trust points
- A trust point is a CA with an advertised policy
- Security domains interlinked by cross-certification among trust points
- User advertises certification path to trust point
- Trust point advertises the cross-certificates it has issued



Personal Security Environment

- Each user securely stores
 - the public key of a trusted user
 - the public key and policy of a trusted CA

Example - two users



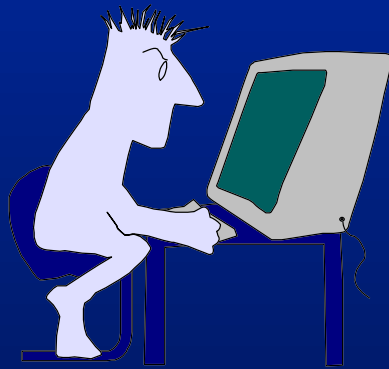
User A



User B

- User A obtains user B's public key by "secure means" and stores it in his PSE.
- User A can authenticate messages from user B
- User B need not do anything
- No policies involved

One user and a small company



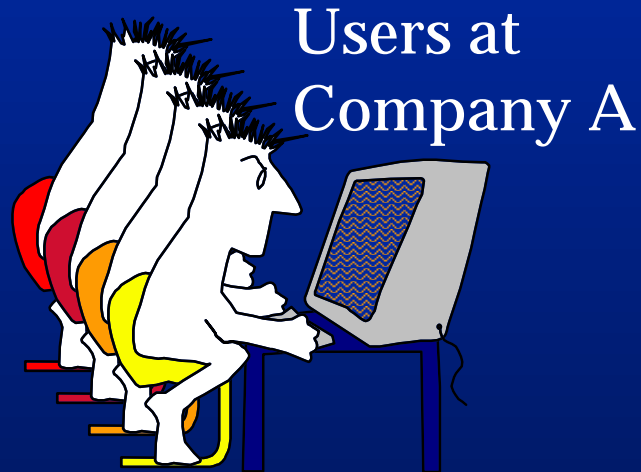
User A



Users at Company B

- Company B creates a CA and publishes a policy
- User A obtains company B's CA's public key and policy and stores it in his PSE.
- User A can authenticate messages from users in company B

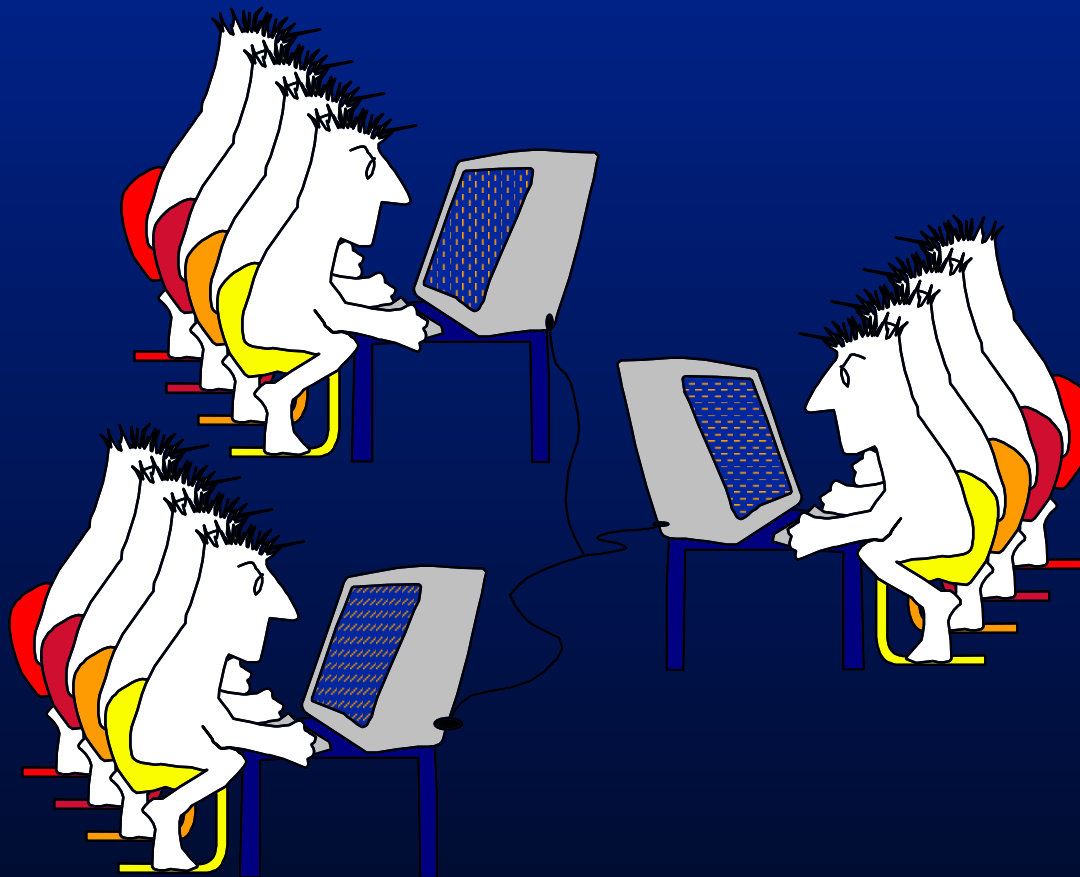
A small company and a big company



- Company B creates a CA hierarchy and publishes a policy for the root CA.
- Company A's CA issues a cross-certificate for Company B's root CA
- Users in company A know their CA's public key and policy.
- Users in company A can authenticate messages from users in company B

Organic growth

Users at Company B





Conclusions

- Scalable deployment model
- Flexibility permits reorganisation
- Supports embedded high security domains
- Explicit use of CA policy

For more information on ICE-TEL

<http://www.darmstadt.gmd.de/ice-tel/ice-home.html>

**Interworking Public Key
Certification Infrastructure for
Europe**



European IVth Framework Project under the TELEMATICS Programme

- 17 partners from 13 countries
- Build and operate CA infrastructure
- Build and pilot secure applications WWW, S/MIME, X.500
- Software from Cost, GMD, Isode, SSE