

Vrije Universiteit Amsterdam

# SafeInit: Comprehensive and Practical Mitigation of Uninitialized Read Vulnerabilities

**Alyssa Milburn**, Herbert Bos,  
Cristiano Giuffrida

# riscure

## Challenge your security

Contact: Alyssa Milburn ([milburn@riscure.com](mailto:milburn@riscure.com))

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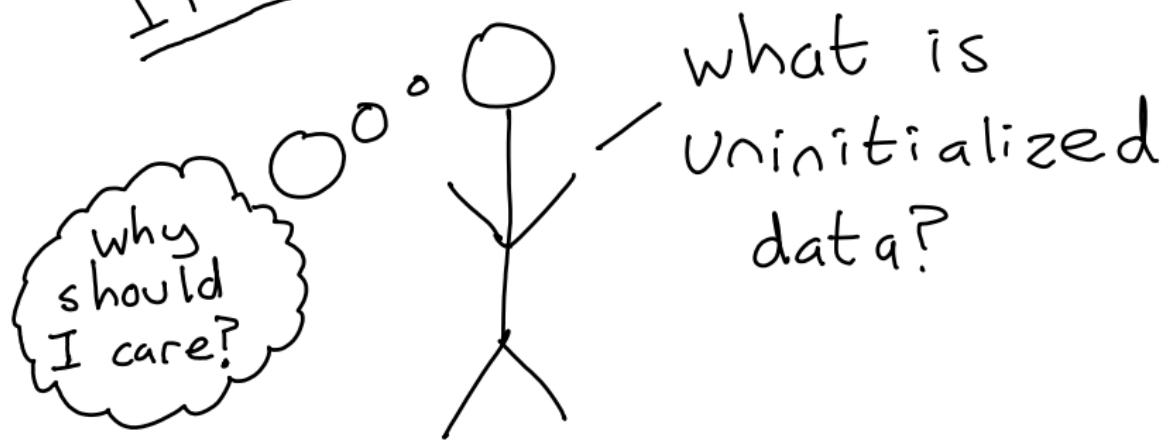
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# INTRODUCTION



# Information exposure

```
keyMemory = malloc();  
*keyMemory = cryptoKeys;
```

Heap allocation is currently used by:  
**keyMemory**

(Secret)  
encryption keys

# Information exposure

```
keyMemory = malloc();  
*keyMemory = cryptoKeys;  
free(keyMemory);
```

Heap allocation is currently used by:  
**nothing**

(Secret)  
encryption keys

# Information exposure

```
keyMemory = malloc();  
*keyMemory = cryptoKeys;  
free(keyMemory);
```

```
buffer = malloc();
```

Heap allocation is currently used by:  
**buffer**

(Secret)  
encryption keys

# Information exposure

```
keyMemory = malloc();  
*keyMemory = cryptoKeys;  
free(keyMemory);
```

```
buffer = malloc();  
send(sock, buffer, ...);
```

Heap allocation is currently used by:

(Secret)  
encryption keys

# Uninitialized read vulnerabilities

# Uninitialized read vulnerabilities

- Information exposure
  - Encryption keys, passwords

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  - Pointers (ASLR)

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- Control flow attacks

# Uninitialized read vulnerabilities

- Information exposure
  - Encryption keys, passwords
  - Pointers (ASLR)
- Control flow attacks
- Undefined behaviour

# Compiler warnings?

```
warning:  
‘variable’ may be used  
uninitialized in this function
```

# Compiler warnings?

github search for -Wno-uninitialized:  
**118,732** code results

But what  
problems  
are you  
trying to  
solve?



qemu

# qemu

CVE-2016-4020: i386: leakage of stack memory to guest in kvmvapic.c

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CVE-2016-4020: i386: leakage of stack memory to guest in kvmvapic.c

```
-     uint32_t imm32;  
+     uint32_t imm32 = 0;
```

# qemu

CVE-2016-4020: i386: leakage of stack memory to guest in kvmvapic.c

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-     uint32_t imm32;  
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CVE-2016-5105: scsi: megasas: stack information leakage while reading configuration

# qemu

CVE-2016-4020: i386: leakage of stack memory to guest in kvmvapic.c

```
-     uint32_t imm32;  
+     uint32_t imm32 = 0;
```

CVE-2016-5105: scsi: megasas: stack information leakage while reading configuration

```
-     uint8_t data[4096];  
+     uint8_t data[4096] = { 0 };
```

# Structure padding



# Structure padding



# Structure padding



# Linux example

# Linux example

CVE-2016-4569: infoleak in Linux sound module

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CVE-2016-4569: infoleak in Linux sound module

```
struct snd_timer_tread tread;  
+ memset(&tread, 0, sizeof(tread));
```

ZERO

MAGIC

ALL

but..

my  
data!



THE  
THINGS

# safeinit

---

# Safelnit

Goal: ensure initialization of the **heap** and **stack**

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- Custom allocator (heap)
- Compiler pass

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```
clang -fsanitize=safeinit  
mycode.c
```

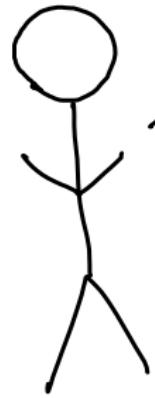
# SafeInit

Goal: ensure initialization of the **heap** and **stack**

- Custom allocator (heap)
- Compiler pass

```
clang -fsanitize=safeinit  
mycode.c
```

- Using clang/LLVM (May 2016) and tcmalloc



But isn't  
that very  
slow?

# Safelnit stack

```
int a;
```

```
a = 1;
```

# Safelnit stack

```
int *a = alloca;
```

```
*a = 1;
```

# Safelnit stack

```
int *a = alloca;
```

```
*a = 0;
```

```
*a = 1;
```

# Safelnit stack

```
int *a = alloca;
```



```
*a = 0;
```

```
*a = 1;
```

# Safelnit stack

```
int *a = alloca;
```

  
Hello,  
\*a = 1;

```
main() {  
    int val; // uninitialized!  
  
    printf("%d", val);  
}
```

# Safelnit stack: IR-level

## baseline

```
define @main() {
    %ptr = alloca i32
    %val = load %ptr
    call printf(..., %val)
}
```

# Safelnit stack: IR-level

## mem2reg

```
define @main() {
    %ptr = alloca i32
    %val = undef
    call printf(..., %val)
}
```

# Safelnit stack: IR-level

## mem2reg

```
define @main() {  
    call printf(..., undef)  
}
```

# Safelnit stack: IR-level

## baseline

```
define @main() {
    %ptr = alloca i32
    %val = load %ptr
    call printf(..., %val)
}
```

# Safelnit stack: IR-level

## Safelnit

```
define @main() {
    %ptr = alloca i32
    call llvm.memset(%ptr, 0, 4) // zero it!

    %val = load %ptr
    call printf(..., %val)
}
```

# Safelnit stack: IR-level

## Safelnit

```
define @main() {
    %ptr = alloca i32
    call llvm.memset(%ptr, 0, 4) // zero it!

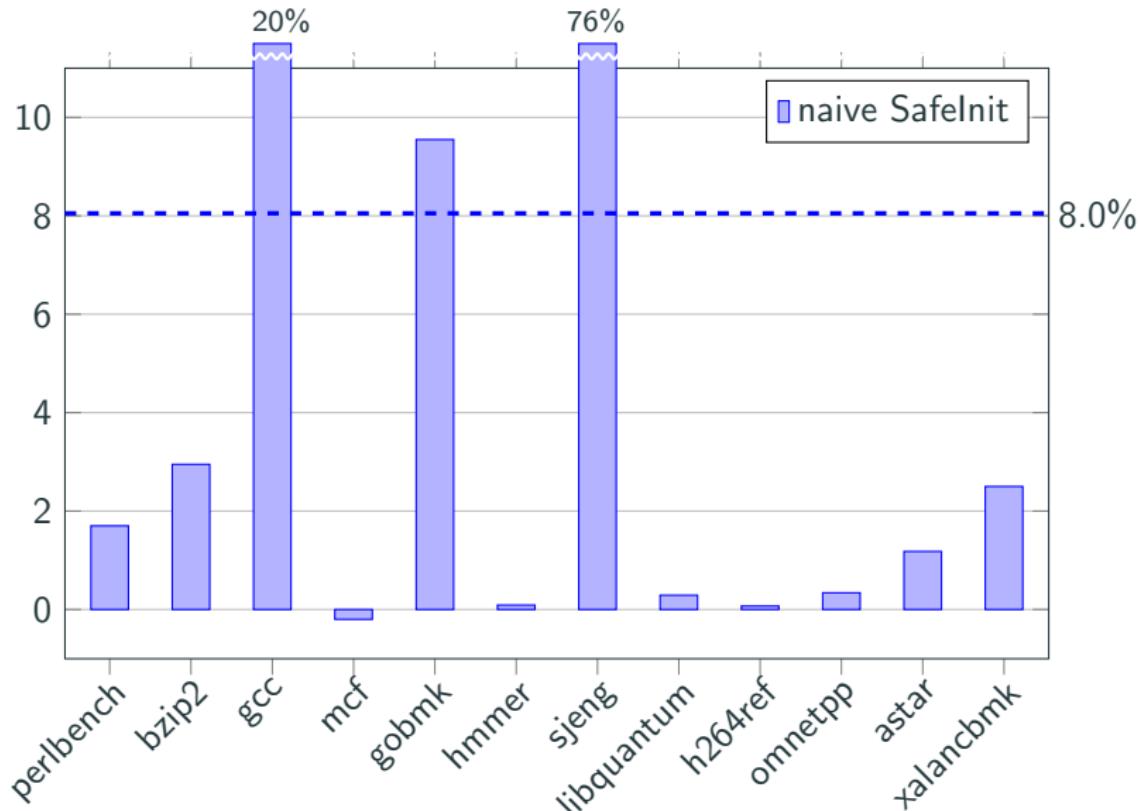
    %val = 0
    call printf(..., %val)
}
```

# Safelnit stack: IR-level

## Safelnit

```
define @main() {  
    call printf(..., 0)  
}
```

# SPEC CINT2006: naive Safelnit



# optimizations

---

# Sinking stores

```
int a, b, c;  
  
...  
  
if (err) {  
    char buf[10000];  
  
    complexPrepare(buf);  
    complexError(buf);  
}
```

# Sinking stores

```
int a, b, c;
```

```
...
```

```
if (err) {
```

```
    char buf[10000];
```

```
    complexPrepare(buf);
```

```
    complexError(buf);
```

```
}
```

 not  
here!

# Sinking stores

```
int a, b, c;
```

```
...
```

```
if (err) {  
    char buf[10000];  
    ←★ here!★  
    complexPrepare(buf);  
    complexError(buf);  
}
```

# More optimizations

- New: Write-only buffer detection

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  - New: Non-constant initialization

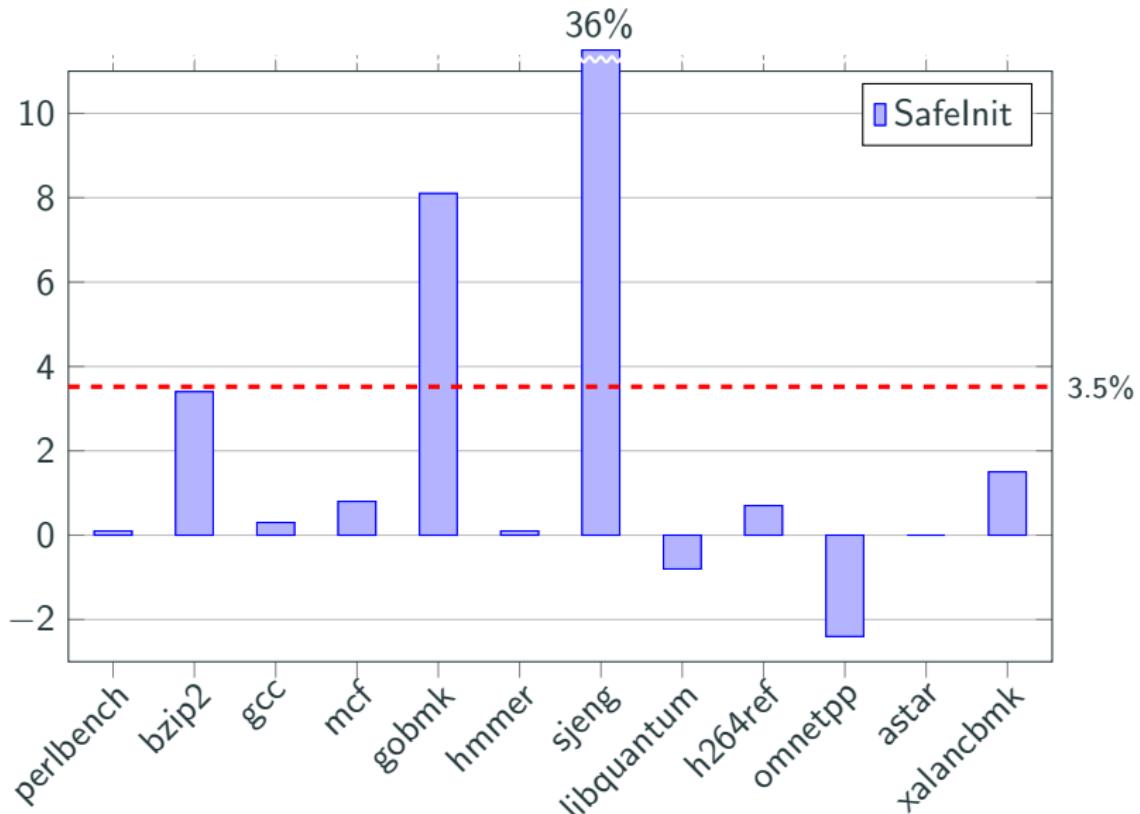
# More optimizations

- New: Write-only buffer detection
- Dead Store Elimination patches
  - New: Non-constant lengths
  - New: Non-constant initialization
- More: ‘Safe’ string buffers, non-constant length store shortening, ...

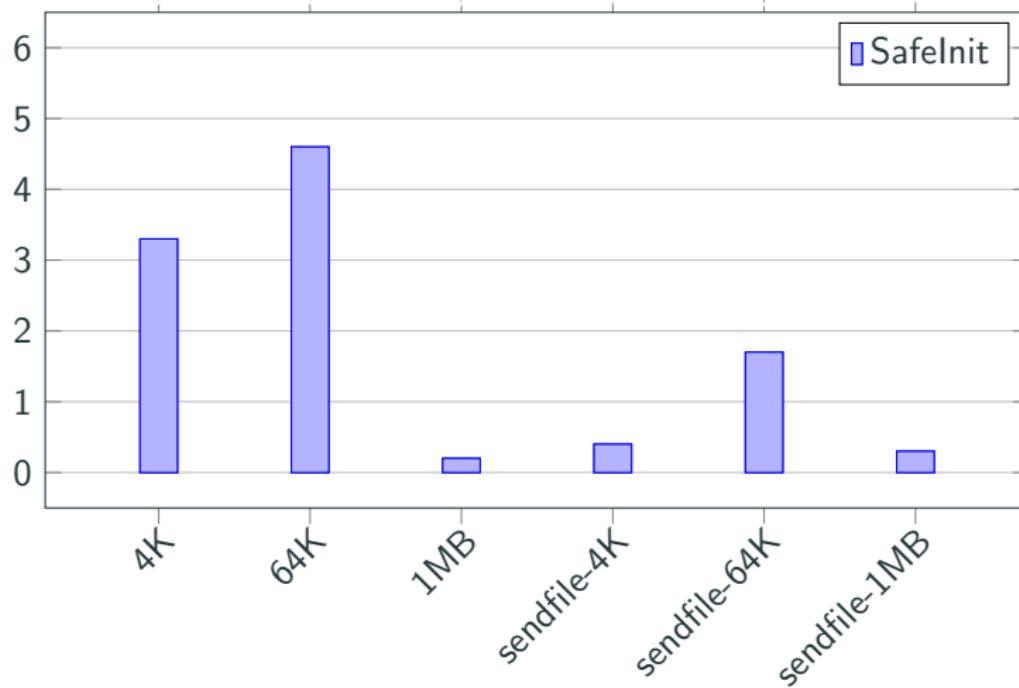
# benchmarks

---

# SPEC CINT2006: SafeInit overhead



# nginx: (worst-case) Safelnit overhead



# Linux: SafelInit overhead

Worst results from LMbench (microbenchmarking system calls) on LLVM Linux:

Sub-benchmark	SafelInit (stack)
no-op system call	0%
fstat system call	4.9%
signal handler (prot fault)	5.9%

# Bonus: actually works

- Often just running valgrind is enough to make it obvious we fixed code!
- Some less trivial CVEs: 2016-4243 (PHP), 2016-5337 (qemu), 2016-4486 (Linux)
- Assembly code sometimes changes drastically!

# Summary

- **SafeInit**: mitigate this entire class of vulnerabilities by simply guaranteeing initialization on stack and heap

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- We obtained acceptable overhead (< 5%)
- Even better as compiler optimizations improve!

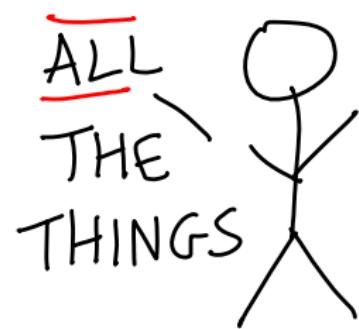
# Summary

- **SafeInit**: mitigate this entire class of vulnerabilities by simply guaranteeing initialization on stack and heap
- We obtained acceptable overhead (< 5%)
- Even better as compiler optimizations improve!
- SafeInit prototype:  
[github.com/vusec/safeinit](https://github.com/vusec/safeinit)

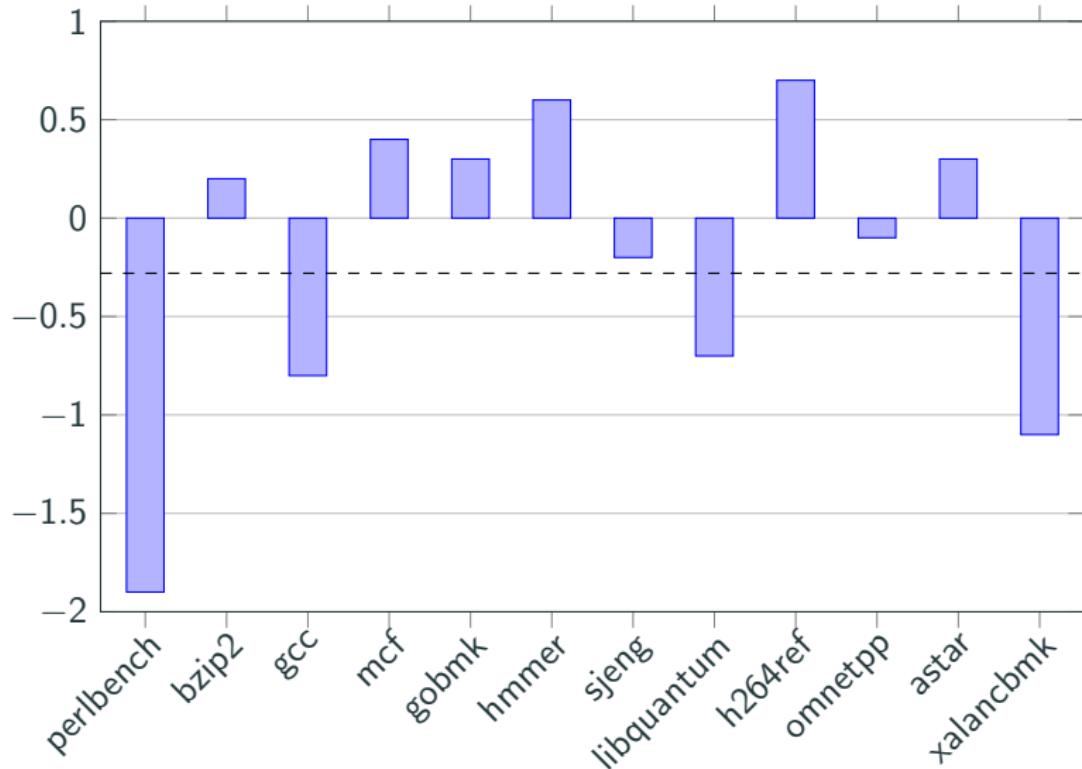
# Summary

- **SafeInit**: mitigate this entire class of vulnerabilities by simply guaranteeing initialization on stack and heap
- We obtained acceptable overhead (< 5%)
- Even better as compiler optimizations improve!
- SafeInit prototype:  
[github.com/vusec/safeinit](https://github.com/vusec/safeinit)
- See the paper for more results and discussion!

# Questions?



# SPEC CINT2006: Optimizer overhead



# Loop initialization

```
int buffer[N];  
memset(buffer, 0, sizeof(buffer));  
  
for (int i = 0; i < N; ++i)  
    buffer[i] = 1;
```

# Loop initialization

```
int buffer[N];  
memset(buffer, 0, sizeof(buffer));  
initialized(buffer, 0, sizeof(buffer));  
for (int i = 0; i < N; ++i)  
    buffer[i] = 1;
```

# Loop initialization

```
int buffer[N];  
memset(buffer, 0, sizeof(buffer));  
initialized(buffer, 0, sizeof(buffer));  
for (int i = 0; i < N; ++i)  
    buffer[i] = 1;
```

# String buffers

```
char buffer[500];
strcpy(buffer, tempString);
strcpy(output, buffer);
```

# Undefined behaviour

```
int deny_access;  
  
if (deny_access) exit();  
  
print_secret_keys();
```

# Undefined behaviour

```
int deny_access = 0;  
  
if (deny_access) exit();  
  
print_secret_keys();
```

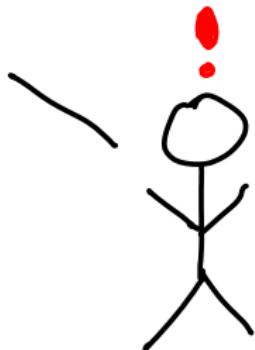
# Undefined behaviour

```
int deny_access = 0;
```

```
if ((deny_access & $) == 0);
```

```
print_secret_keys();
```

Oh-oh...



# Annotations

Huge zero initialization getting you down?

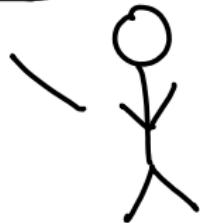
`__attribute__((no_zeroinit))`

# Annotations

Huge zero initialization getting you down?

`__attribute__((no_zeroinit))`

but you said automated!

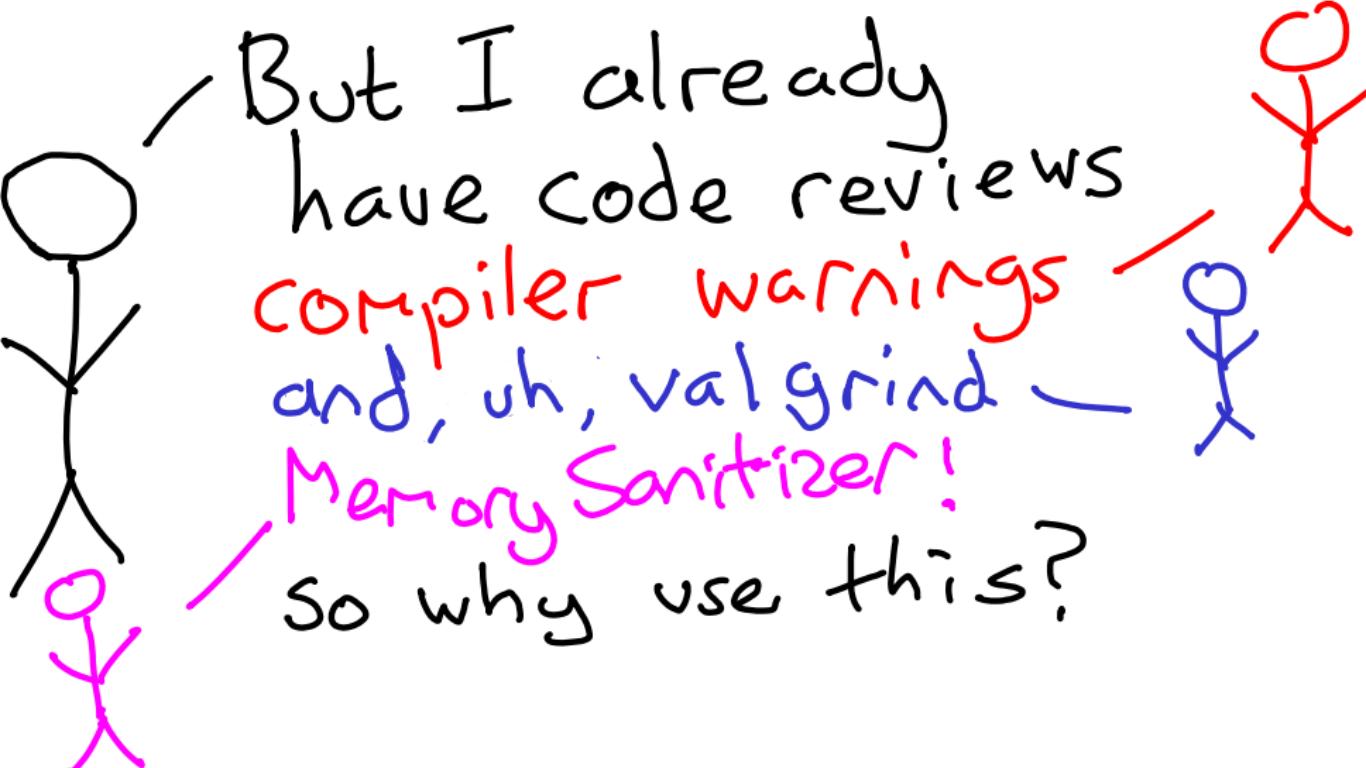


# Annotations

Huge zero initialization getting you down?

```
__attribute__((no_zeroinit))
```

WARNING: Excessive size stack allocation of type  
move\_s[500] in test.c on line 20



But I already  
have code reviews

compiler warnings

and, uh, valgrind —

Memory Sanitizer!

So why use this?

# Related Work: mitigation

Heap allocators

Debug allocators, jemalloc, ...

# Related Work: mitigation

## Secure deallocation

“Shredding your garbage”

by Chow et al.

Clear heap memory in free!

# Related Work: mitigation

## Secure deallocation

Frame clearing in epilogue:

“10% – 40%” runtime overhead

(we tried clearing in prologue:  
still > 10%)

# Related Work: mitigation

## PaX gcc plugins

- Stackleak
- Structleak

## Related Work: mitigation

UniSan

Lu et al, CCS 2016.

Kernel info exposure:  
static analysis + initialization