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Joint work with Gian Marti, Jan Müller and Laurent Vanbever

Partition Attack

An adversary splits the Bitcoin network in two disjoint components



Any Blockchain system is vulnerable

Any Blockchain system is vulnerable

Double-spending, Revenue Loss, DoS

Any Blockchain system is vulnerable

Double-spending, Revenue Loss, DoS

50-50 partition is feasible

Any Blockchain system is vulnerable

Double-spending, Revenue Loss, DoS

50-50 partition is feasible

Any network in the world is a possible attacker

In 2017 we uncovered the practicality and effectiveness of routing attacks in Bitcoin

Hijacking Bitcoin: Routing Attacks on Cryptocurrencies

https://btc-hijack.ethz.ch

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Abstract—As the most successful cryptocurrency to date, Bitcoin constitutes a target of choice for attackers. While many attack vectors have already been uncovered, one important vector has been left out though: attacking the currency via the Internet routing infrastructure itself. Indeed, by manipulating routing advertisements (BGP hijacks) or by naturally intercepting traffic, Autonomous Systems (ASes) can intercept and manipulate a large fraction of Bitcoin traffic.

This paper presents the first taxonomy of routing attacks and their impact on Bitcoin, considering both small-scale attacks, targeting individual nodes, and large-scale attacks, targeting the network as a whole. While challenging, we show that two key properties make routing attacks practical: (*i*) the efficiency of routing manipulation; and (*ii*) the significant centralization of Bitcoin in terms of mining and routing. Specifically, we find that any network attacker can hijack few (<100) BGP prefixes to isolate ~50% of the mining power—even when considering that mining pools are heavily multi-homed. We also show that on-path network attackers can considerably slow down block propagation by interfering with few key Bitcoin messages.

We demonstrate the feasibility of each attack against the deployed Bitcoin software. We also quantify their effectiveness on the current Bitcoin topology using data collected from a Bitcoin supernode combined with BGP routing data.

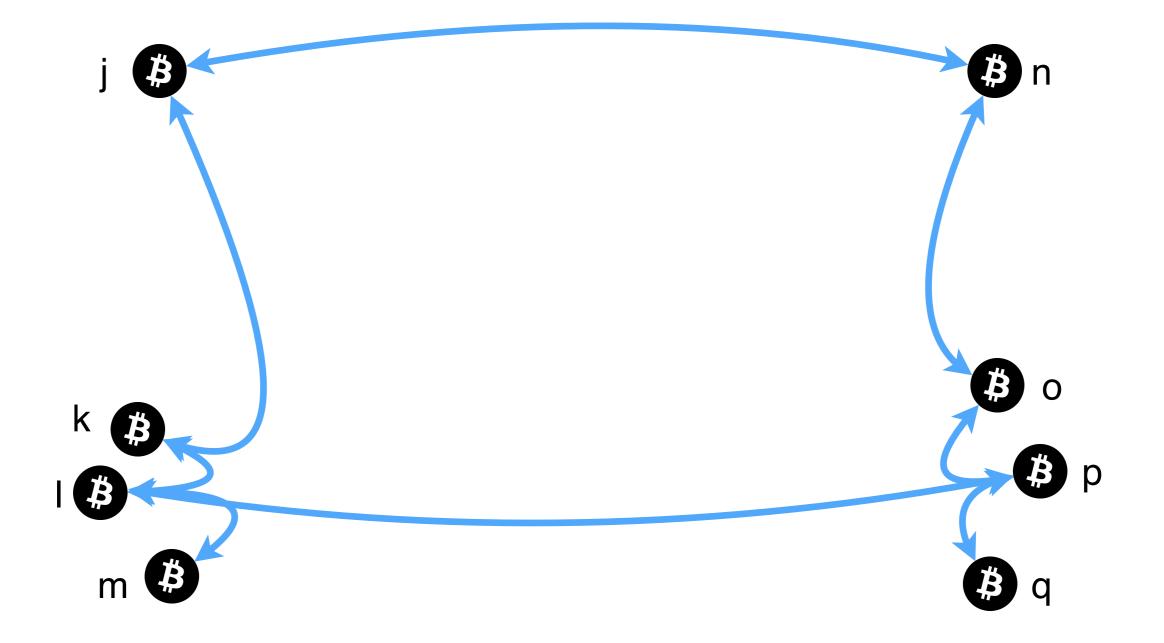
The potential damage to Bitcoin is worrying. By isolating parts of the network or delaying block propagation, attackers can cause

One important attack vector has been overlooked though: attacking Bitcoin via the Internet infrastructure using routing attacks. As Bitcoin connections are routed over the Internet in clear text and without integrity checks-any third-party on the forwarding path can eavesdrop, drop, modify, inject, or delay Bitcoin messages such as blocks or transactions. Detecting such attackers is challenging as it requires inferring the exact forwarding paths taken by the Bitcoin traffic using measurements (e.g., traceroute) or routing data (BGP announcements), both of which can be forged [41]. Even ignoring detectability, mitigating network attacks is also hard as it is essentially a human-driven process consisting of filtering, routing around or disconnecting the attacker. As an illustration, it took Youtube close to 3 hours to locate and resolve rogue BGP announcements targeting its infrastructure in 2008 [6]. More recent examples of routing attacks such as [51] (resp. [52]) took 9 (resp. 2) hours to resolve in November (resp. June) 2015.

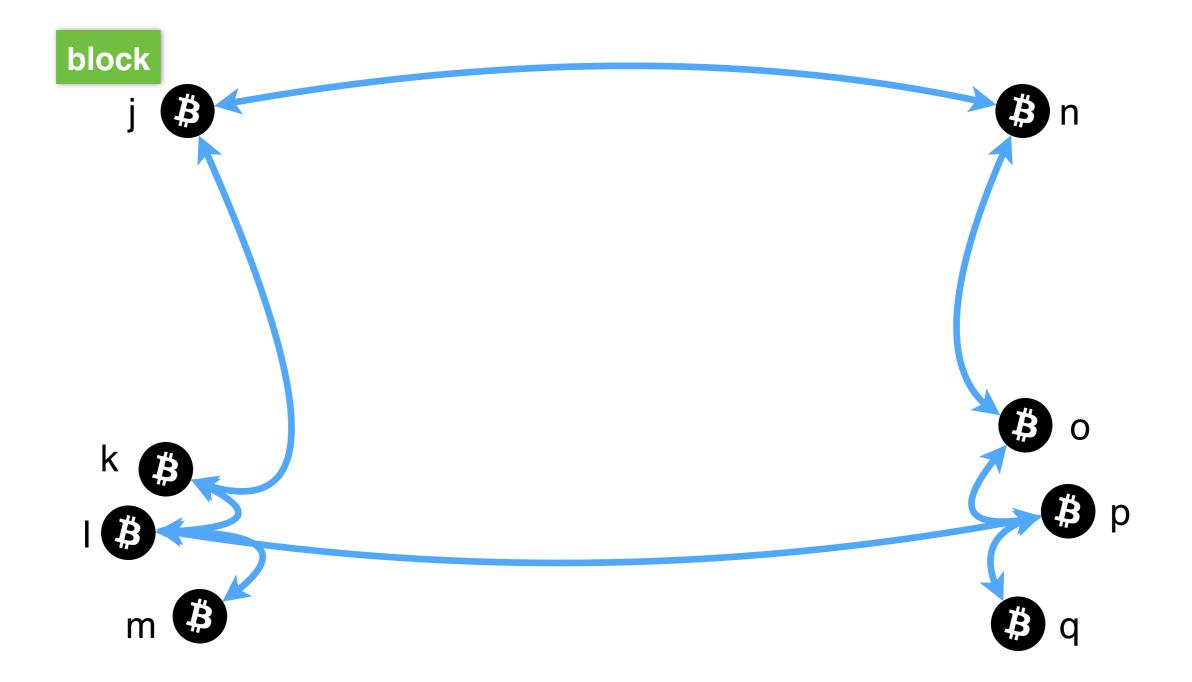
One of the reasons why routing attacks have been overlooked in Bitcoin is that they are often considered too challenging to be practical. Indeed, perturbing a vast peer-to-peer Bitcoin is a distributed network of nodes (Bitcoin clients)



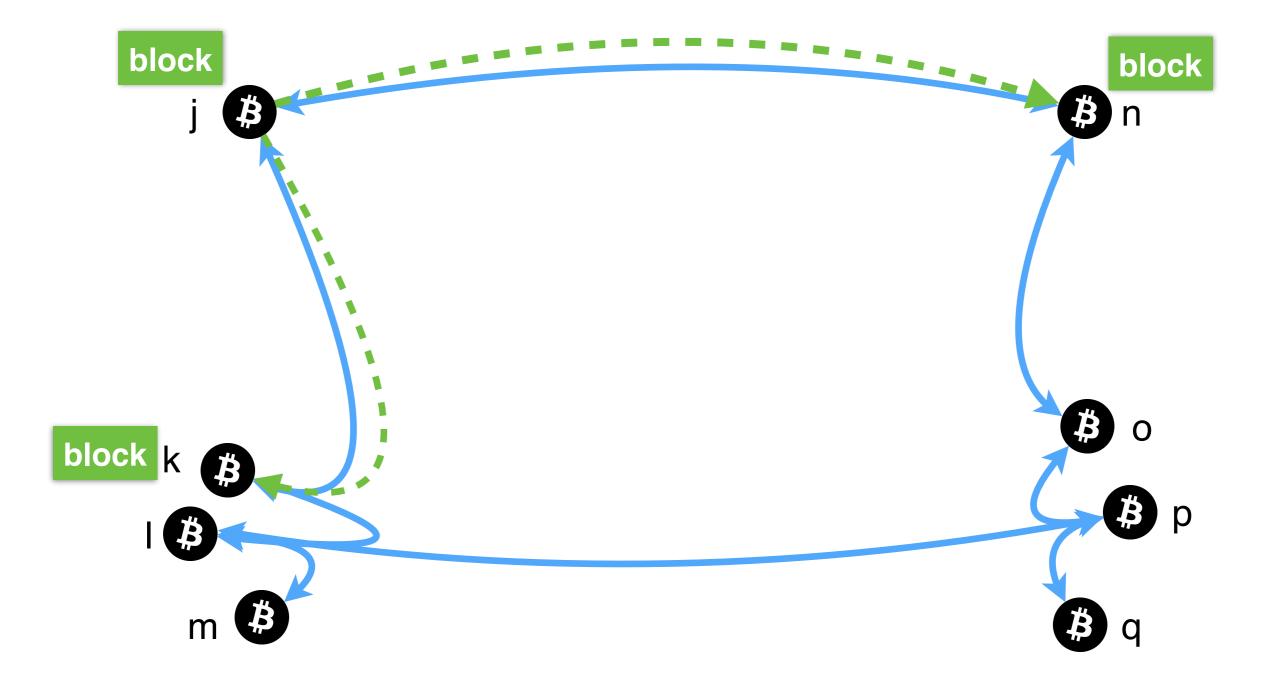
Bitcoin clients establish random connections



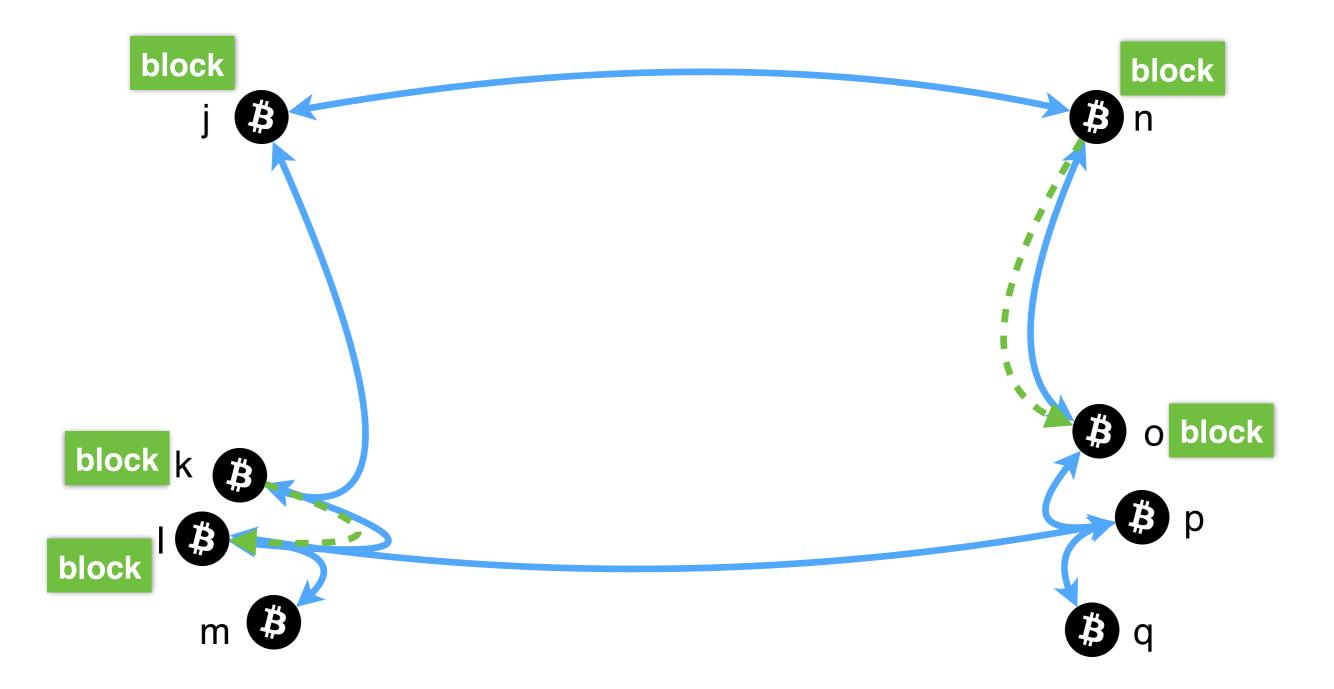
Bitcoin clients exchange Blocks



Blocks contain the latest transactions

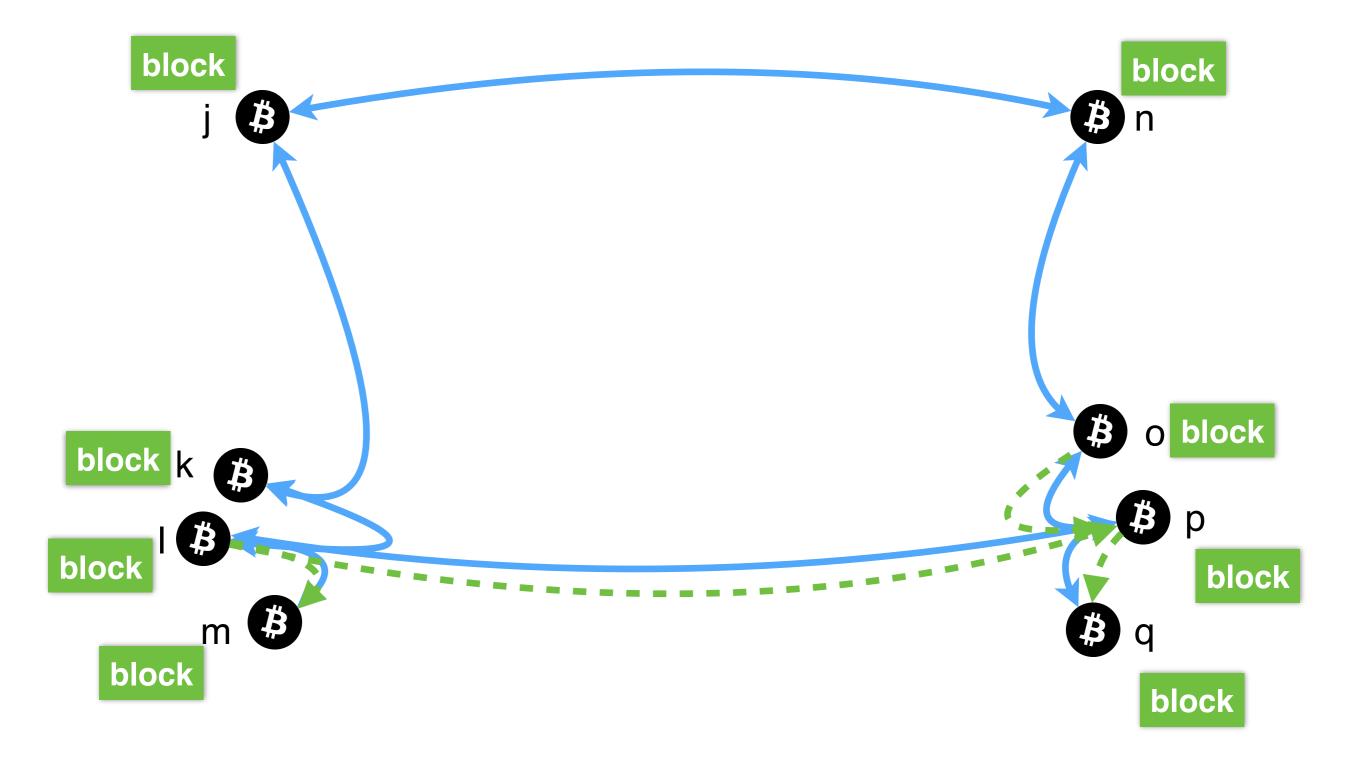


Bitcoin clients exchange Blocks



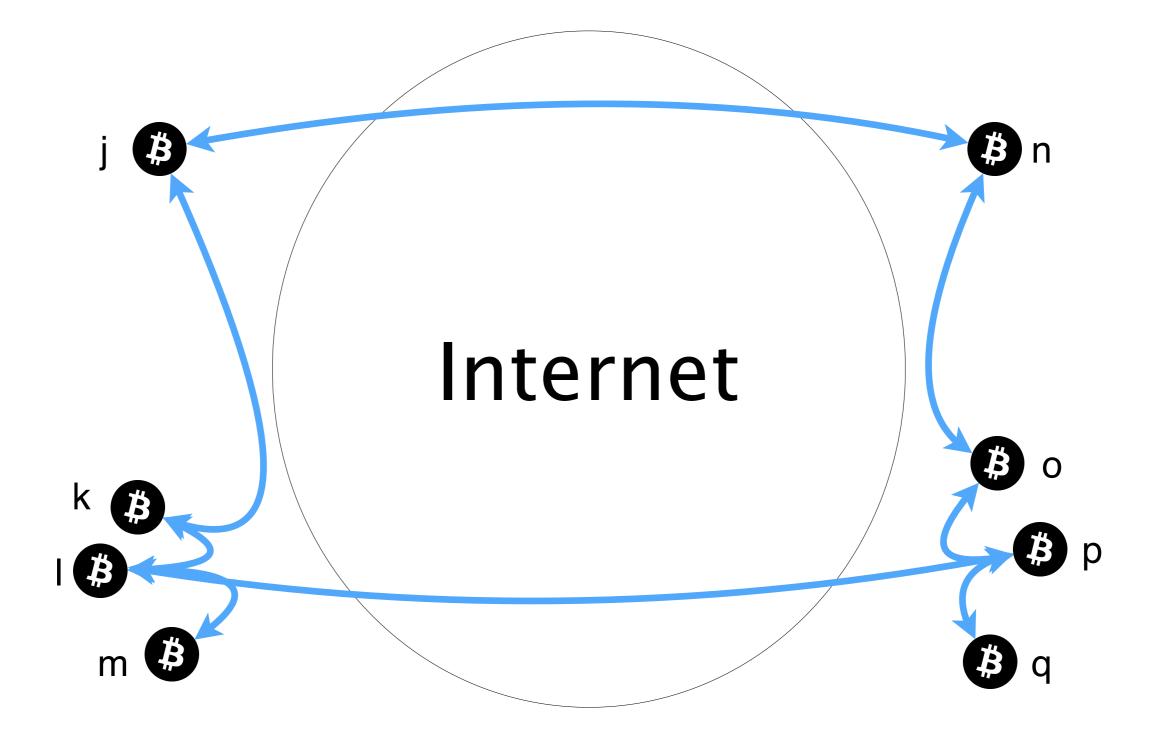
Bitcoin clients exchange Blocks

until all clients have the same view of the transactions

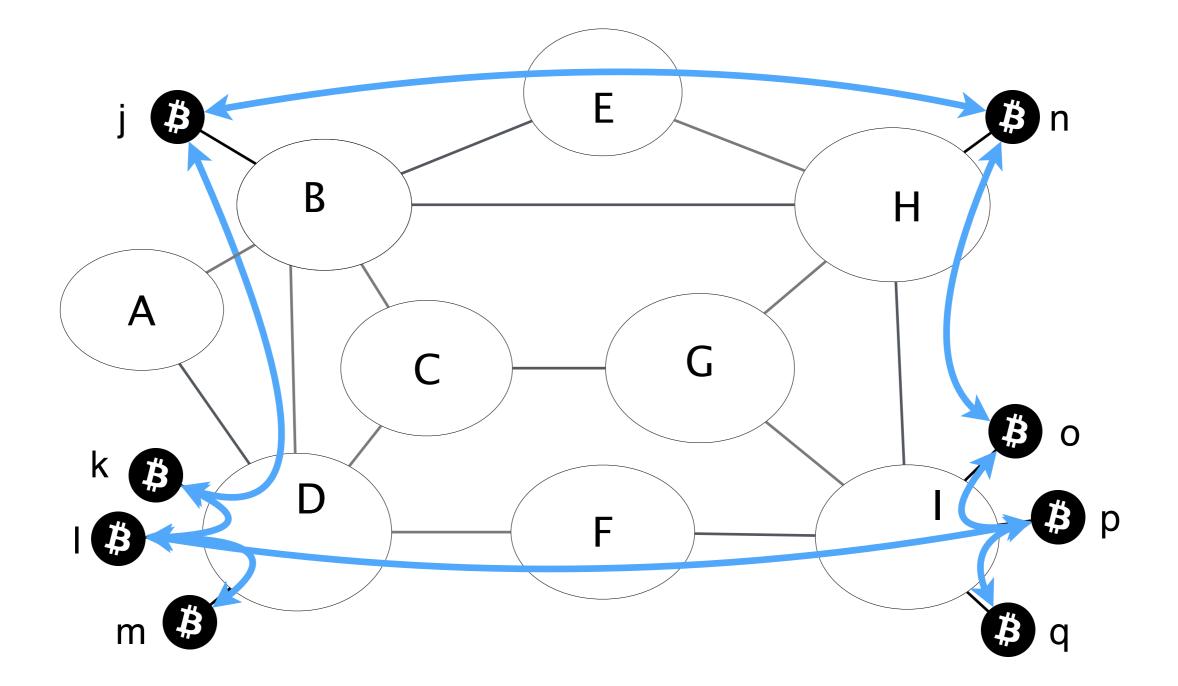


What can go wrong?

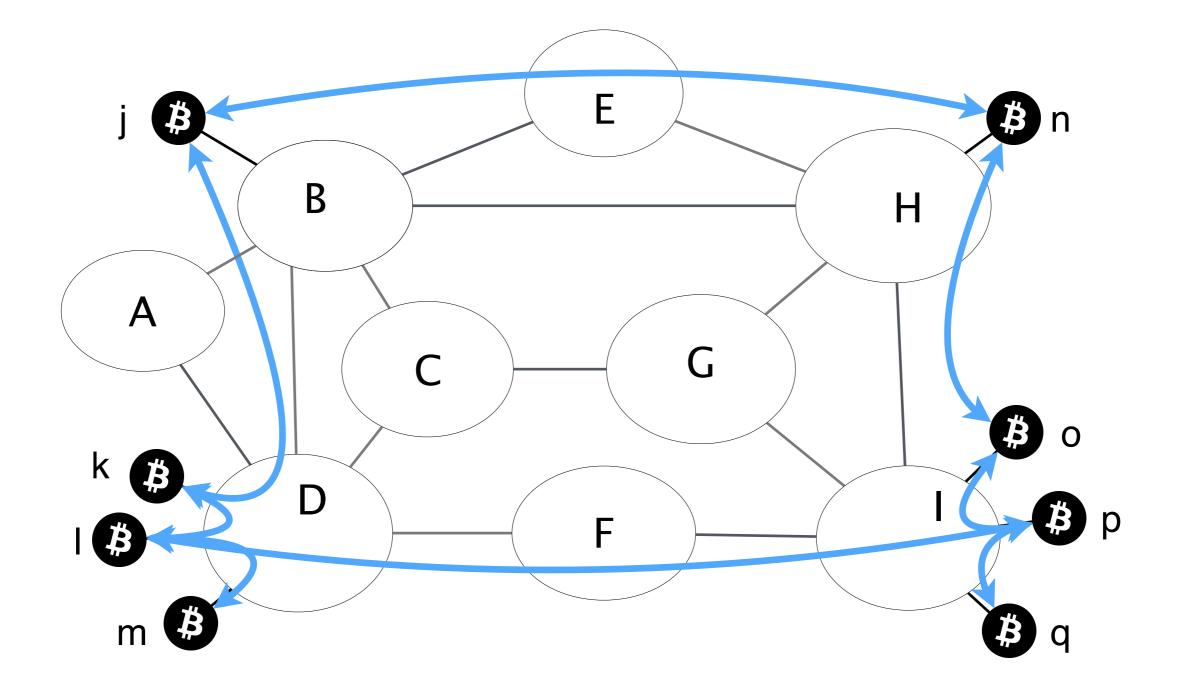
Bitcoin connections are routed over the Internet



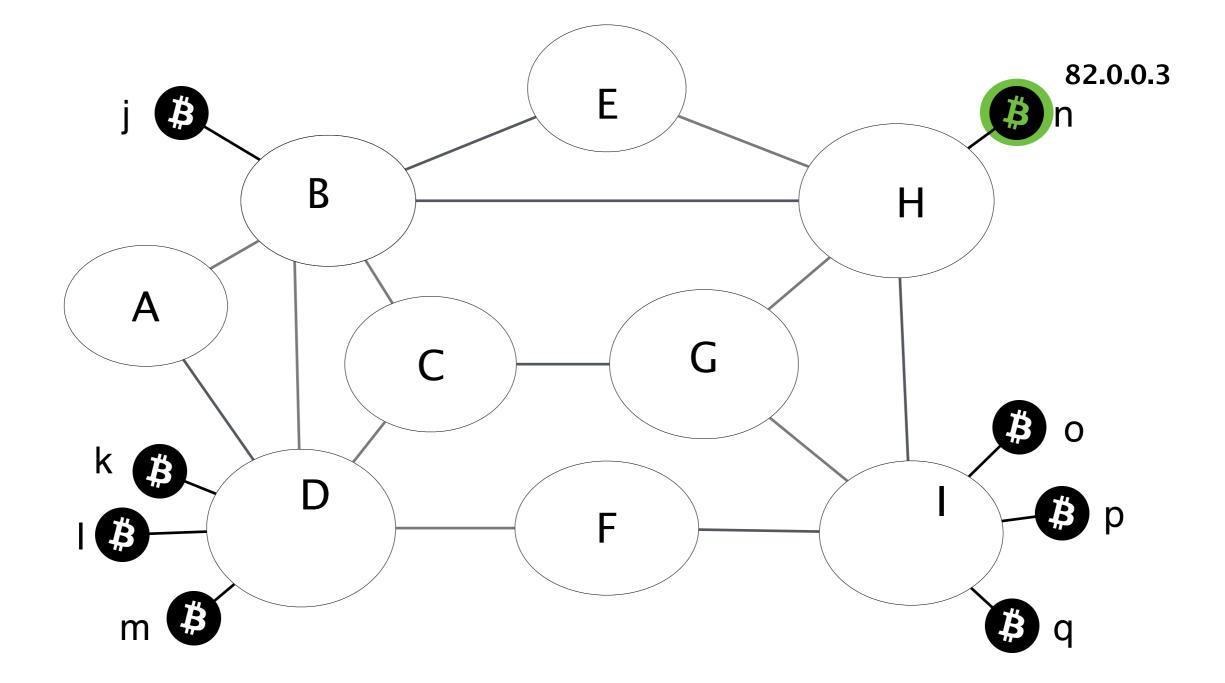
The Internet is composed of Autonomous Systems



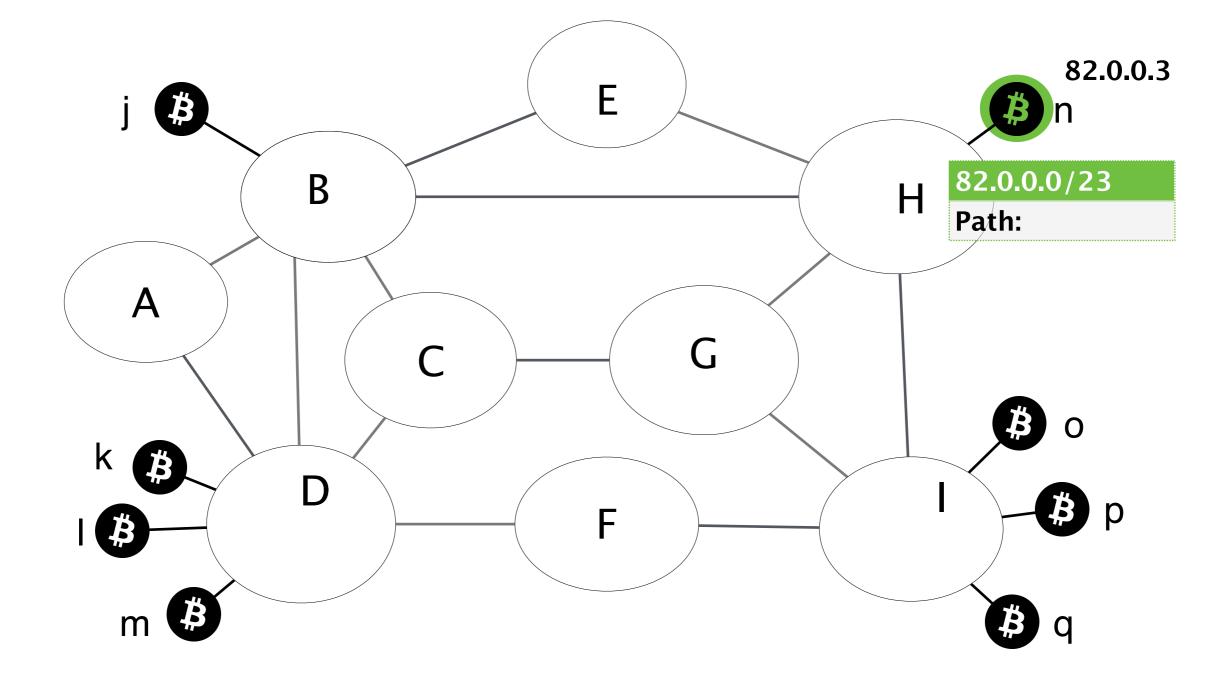
BGP is the default Internet routing protocol



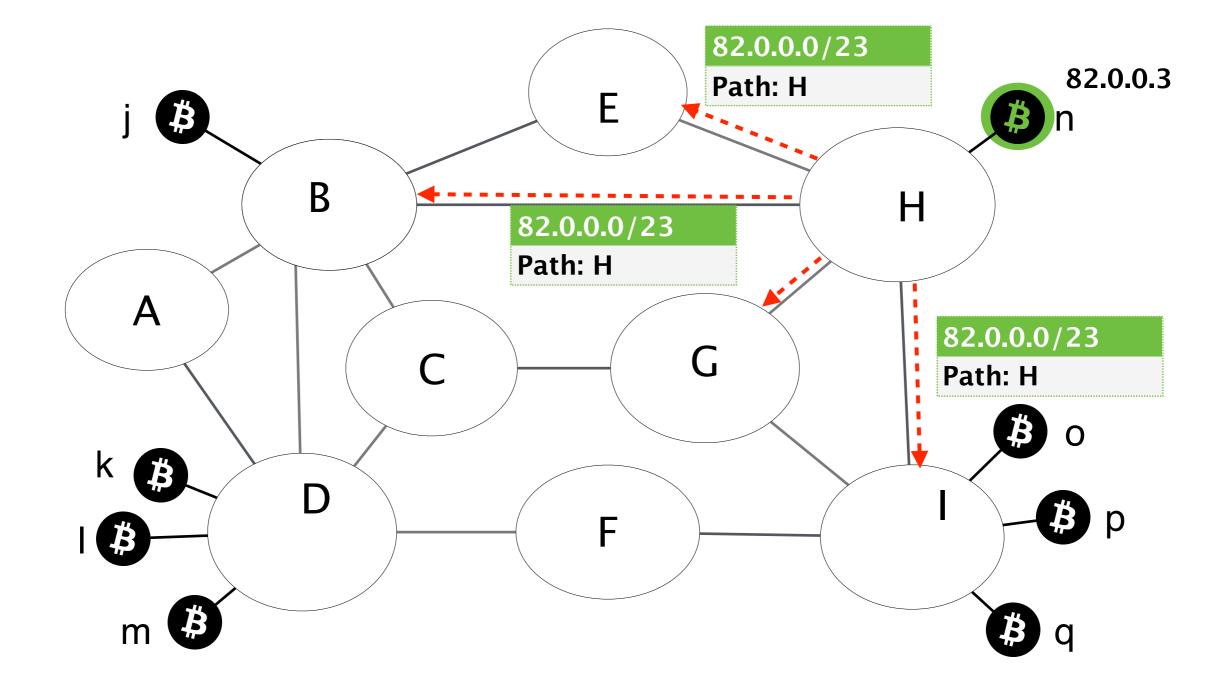
Each Bitcoin client n has an IP



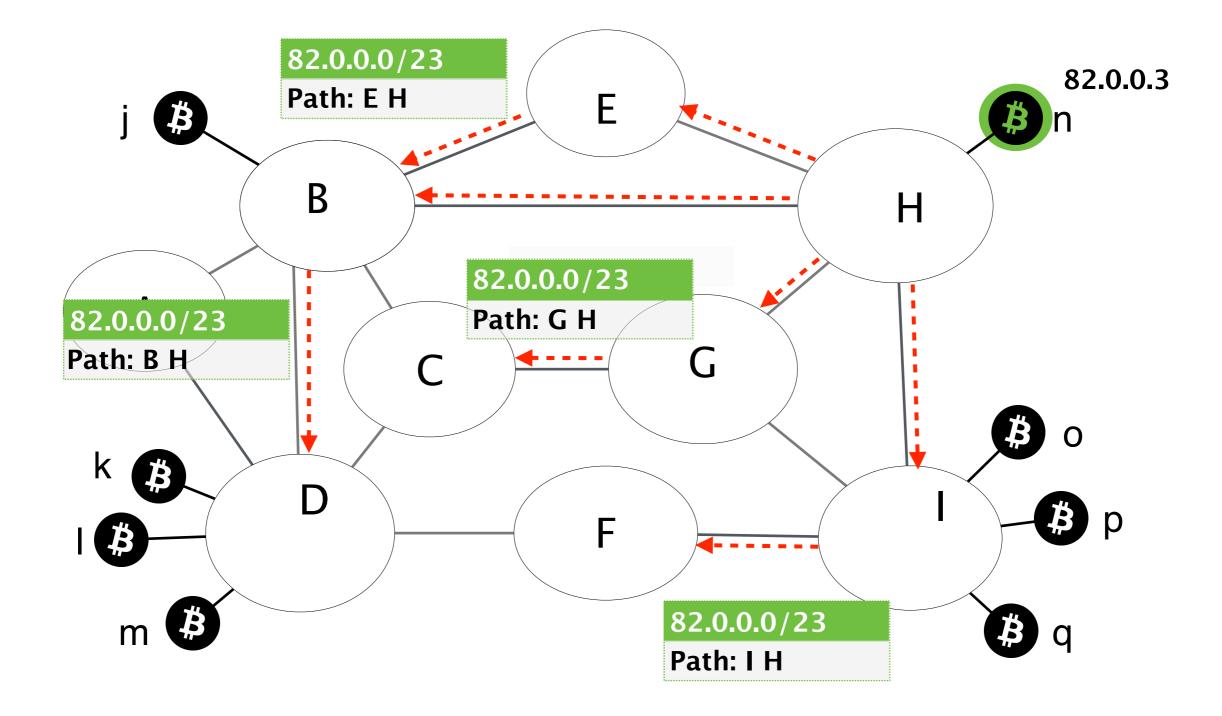
AS H creates a BGP advertisement for n's IP prefix



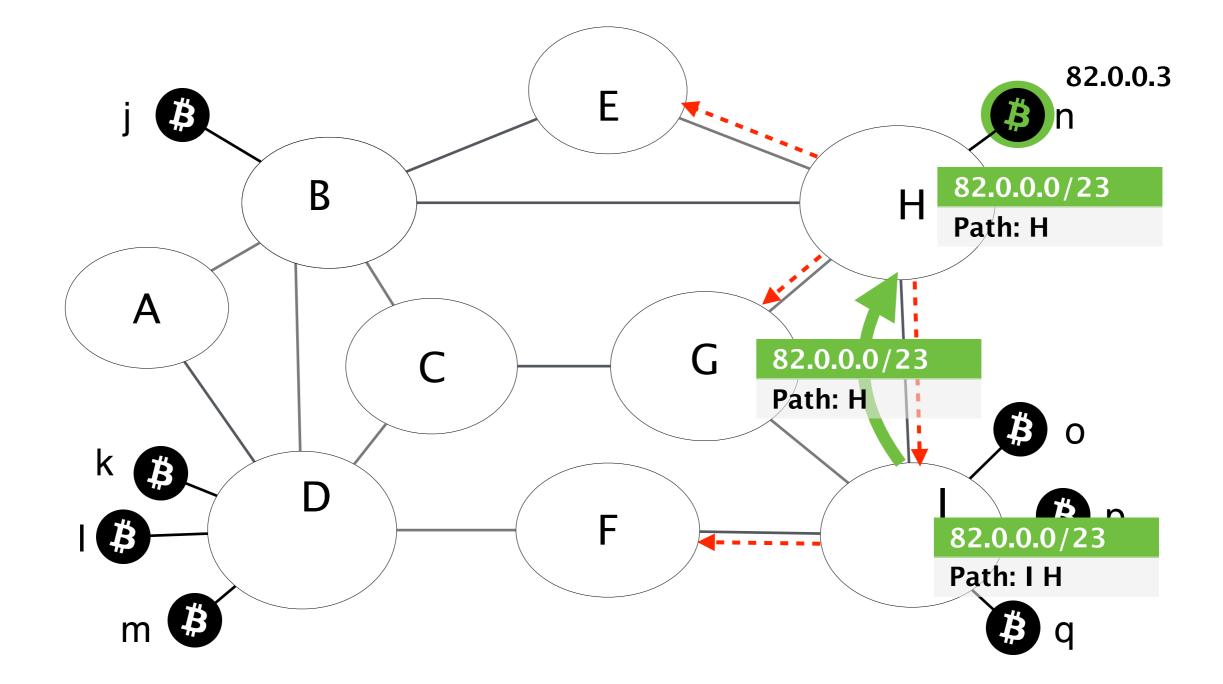
BGP propagates advertisements in the Internet



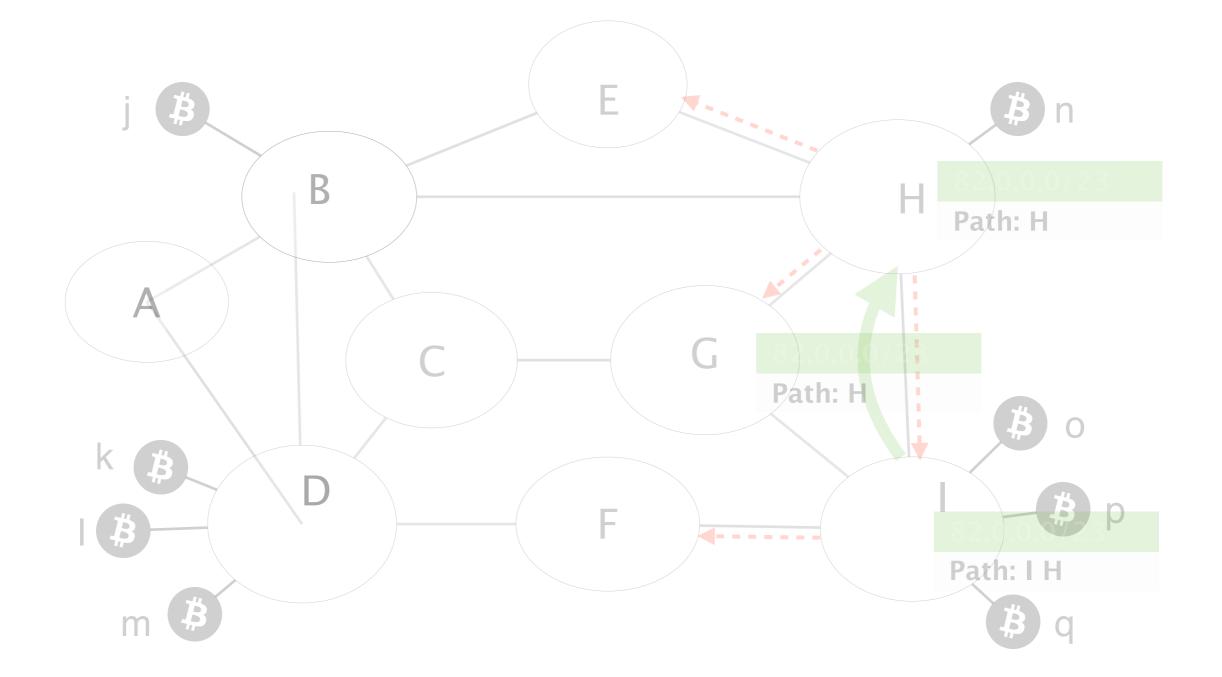
BGP propagates advertisements in the Internet



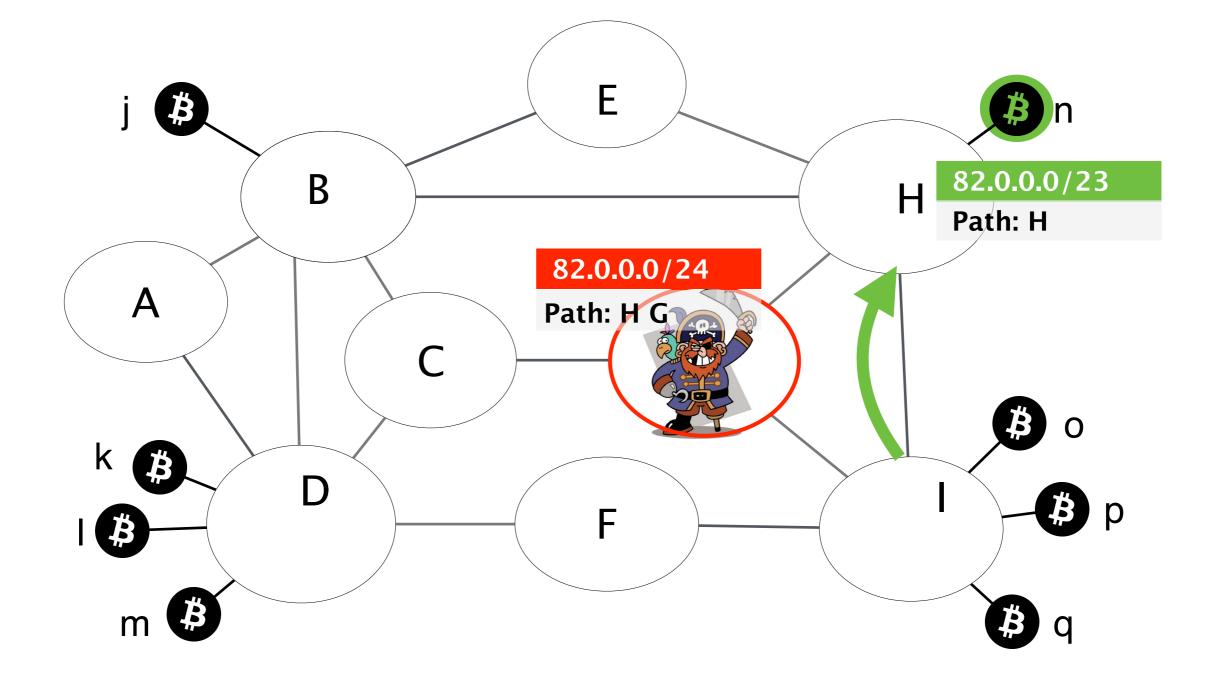
AS I can directly reach AS H



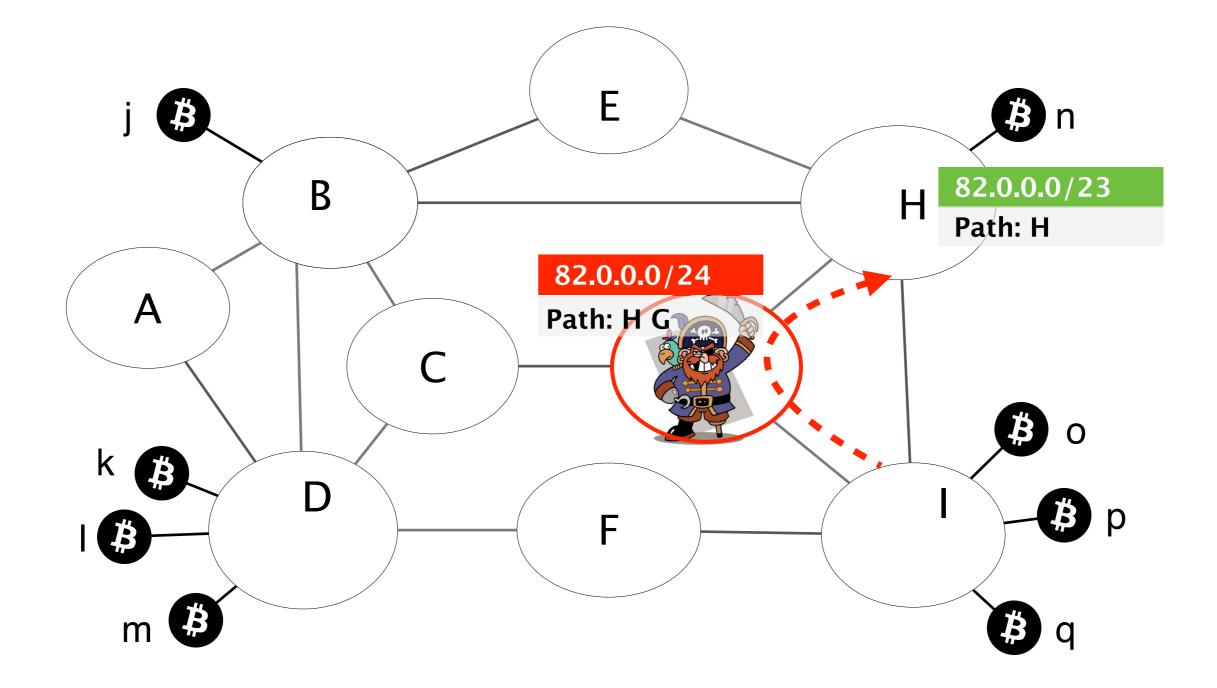
BGP does not check the legitimacy of advertisements



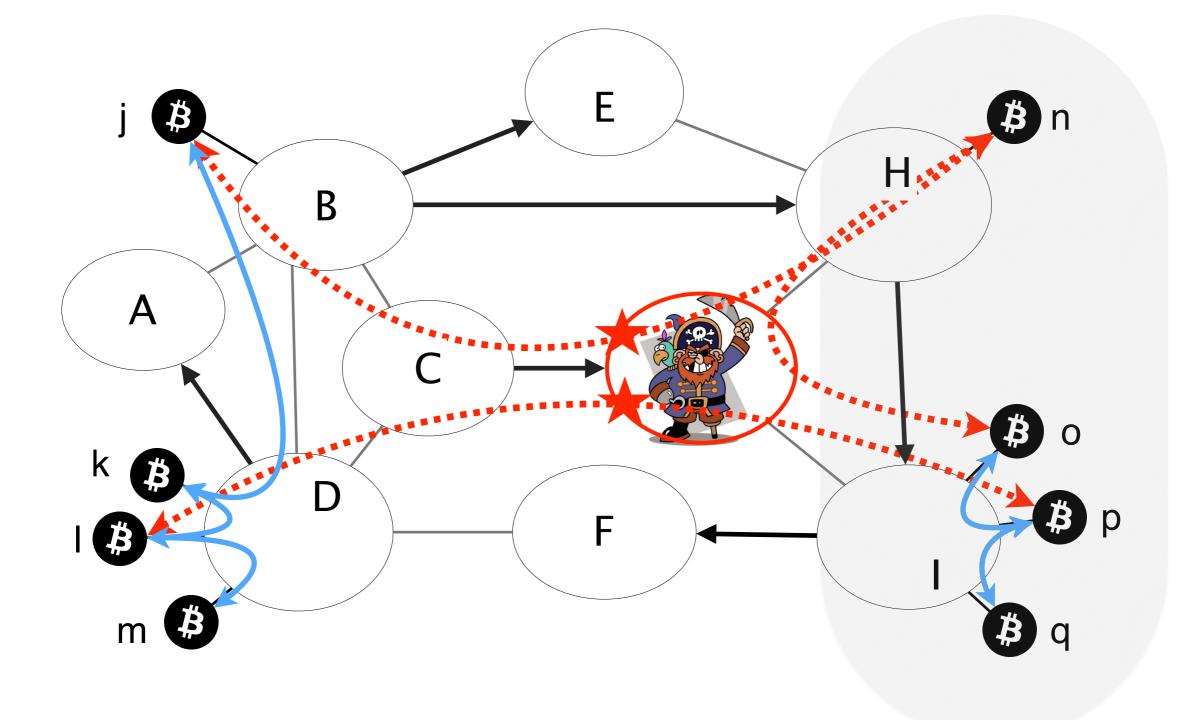
Attacker creates a fake BGP advertisement



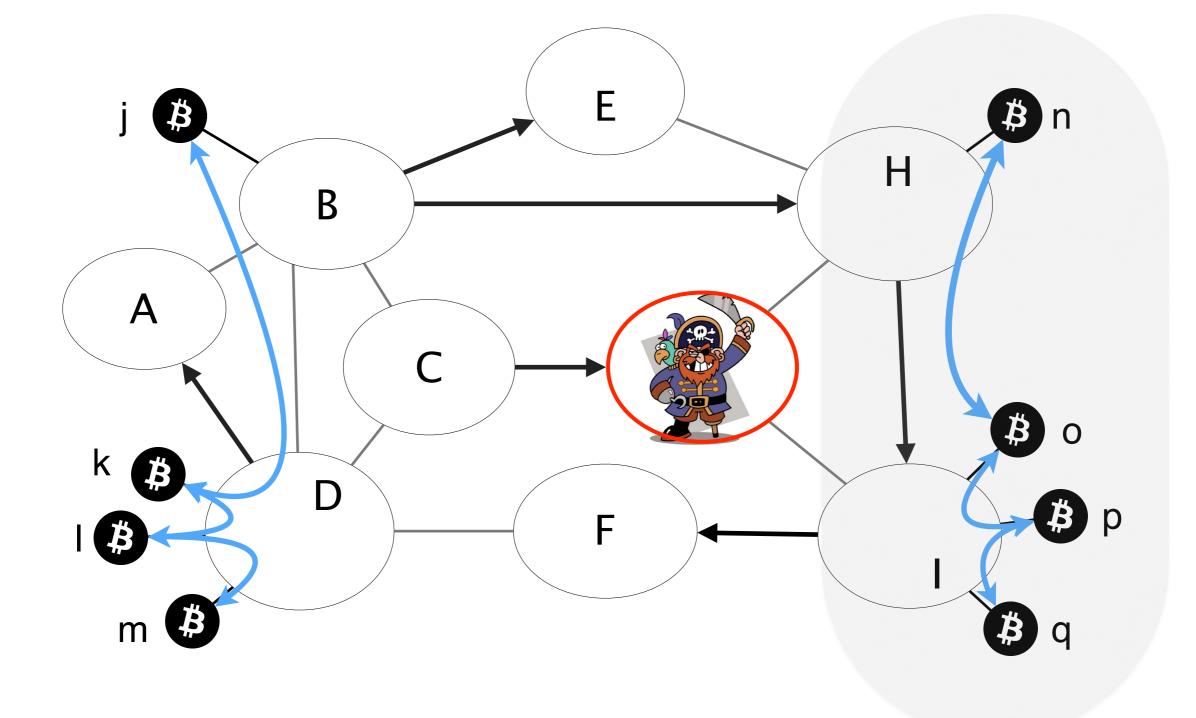
Attacker attracts traffic destined to AS H using BGP hijacking



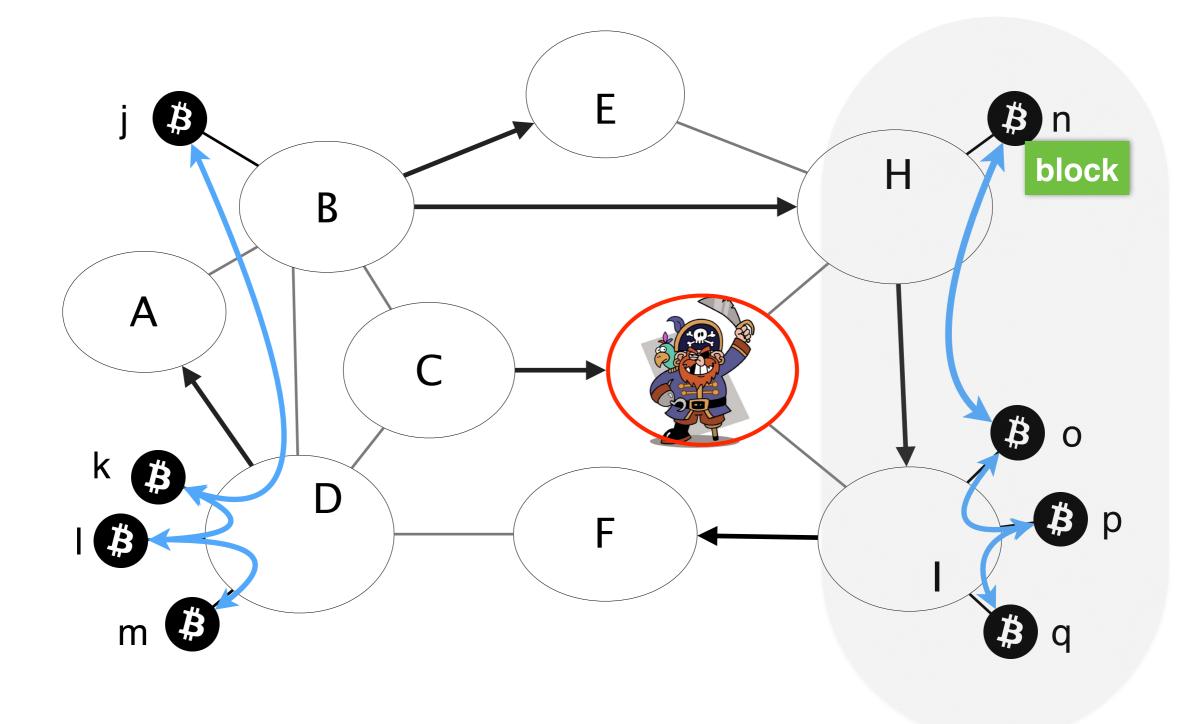
Attacker attracts connections with BGP hijacking



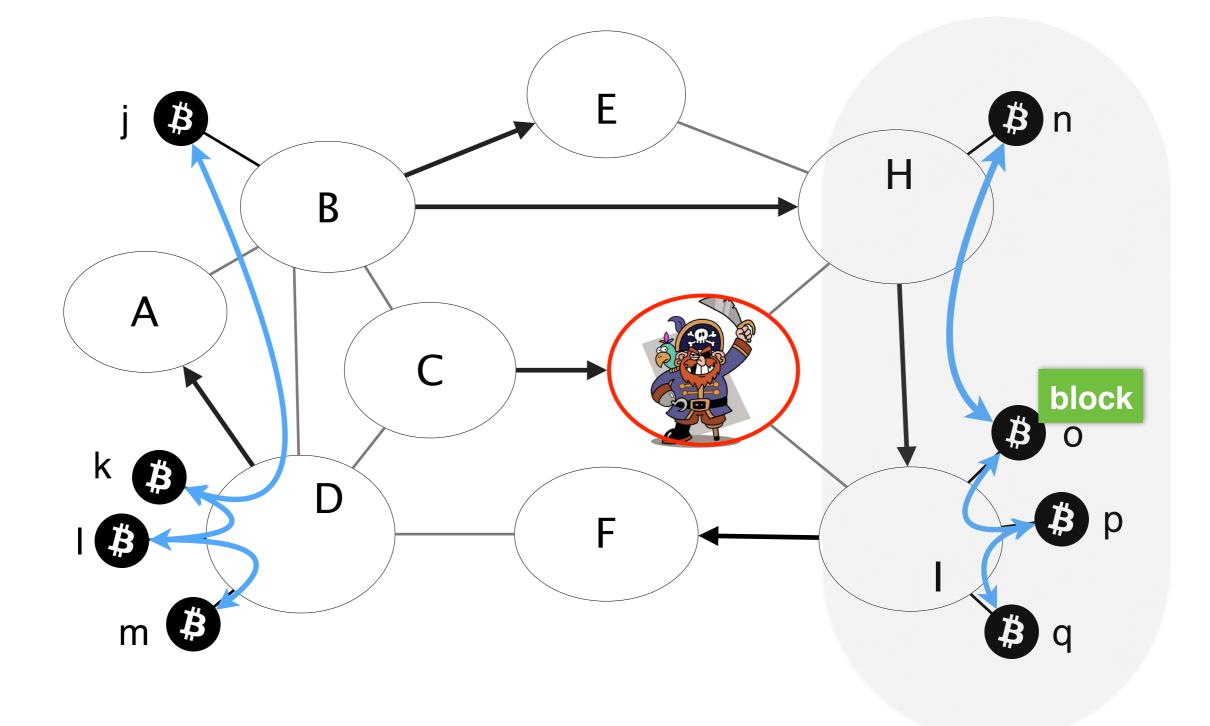
Attacker drops connections crossing the partition



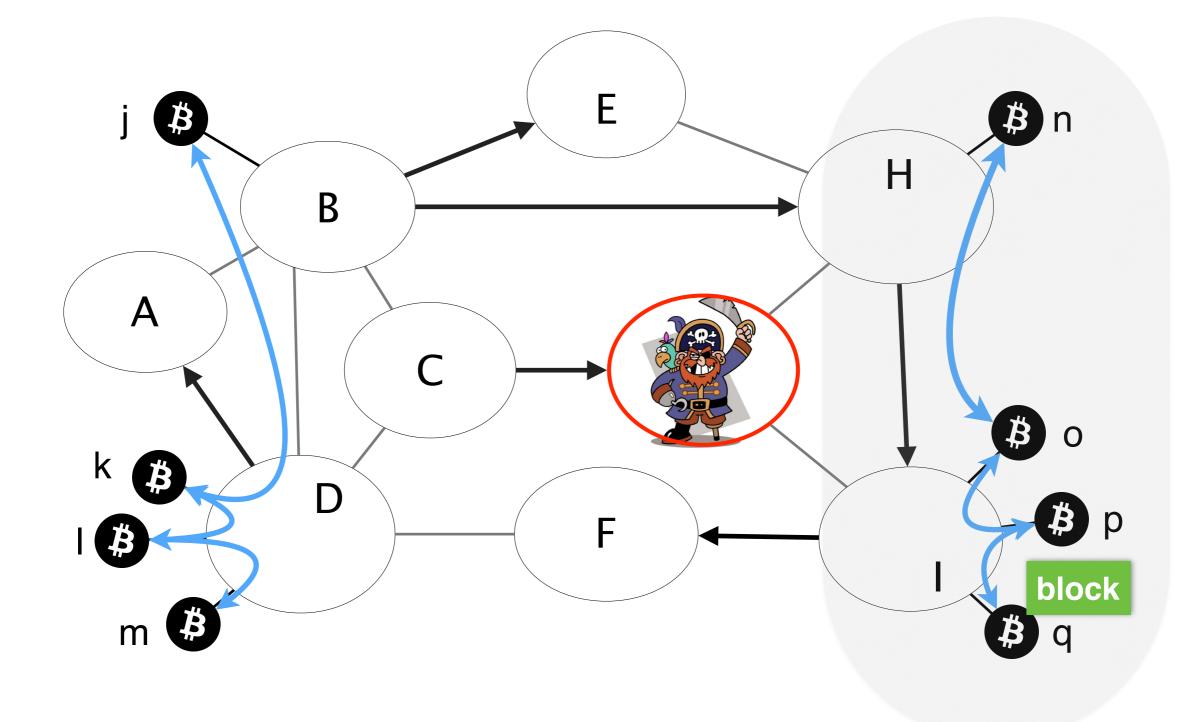
A new block in the grey zone cannot be propagated further



A new block in the grey zone cannot be propagated further



A new block in the grey zone cannot be propagated further



SABRE:

Additional overlay network that is engineered to allow clients to exchange blocks, even if the Bitcoin network is partitioned



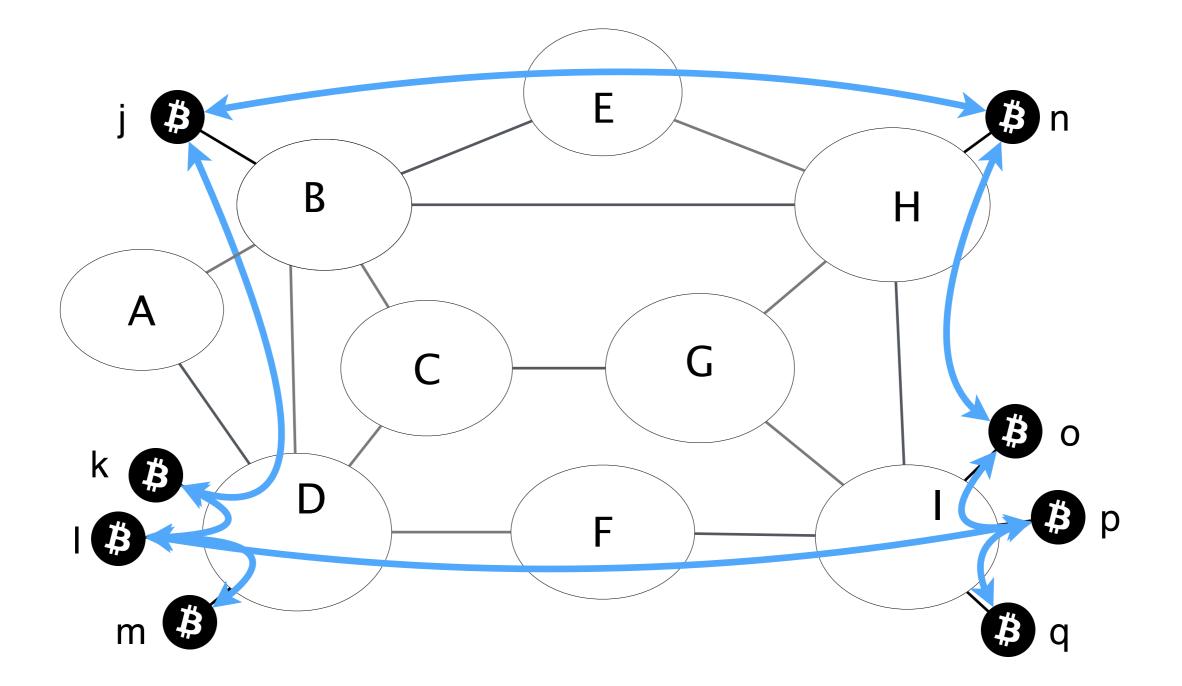
SABRE:

Additional overlay network that is engineered to allow clients to exchange blocks, even if the Bitcoin network is partitioned

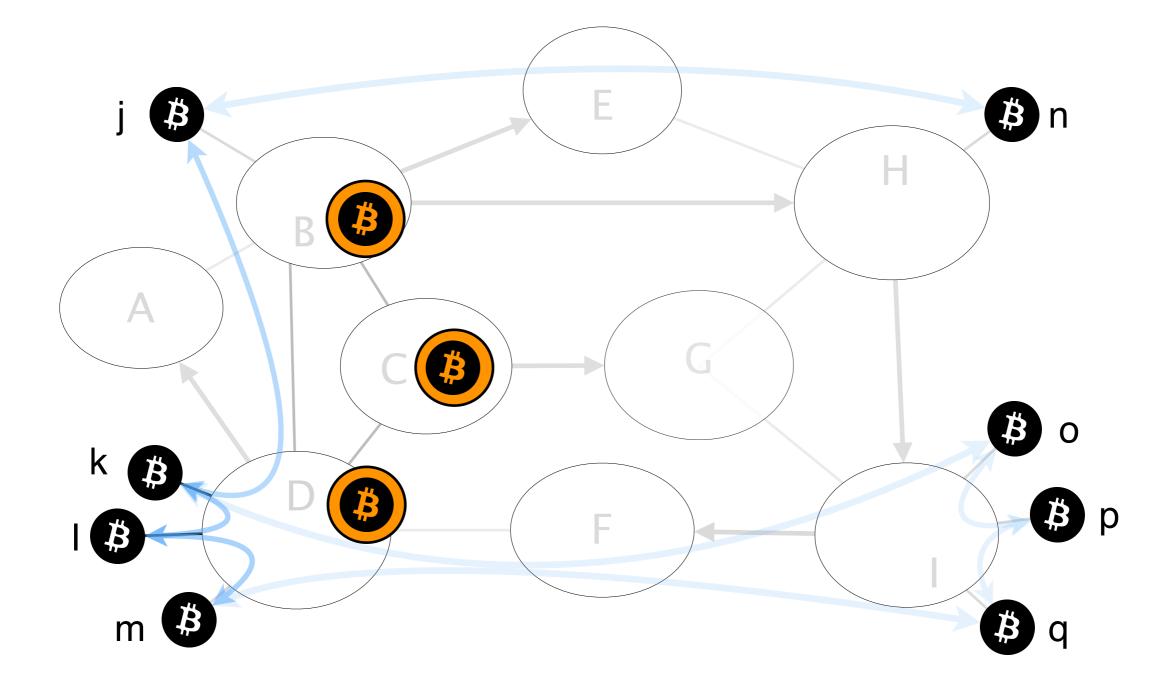
... without the need to deploy secure routing protocols



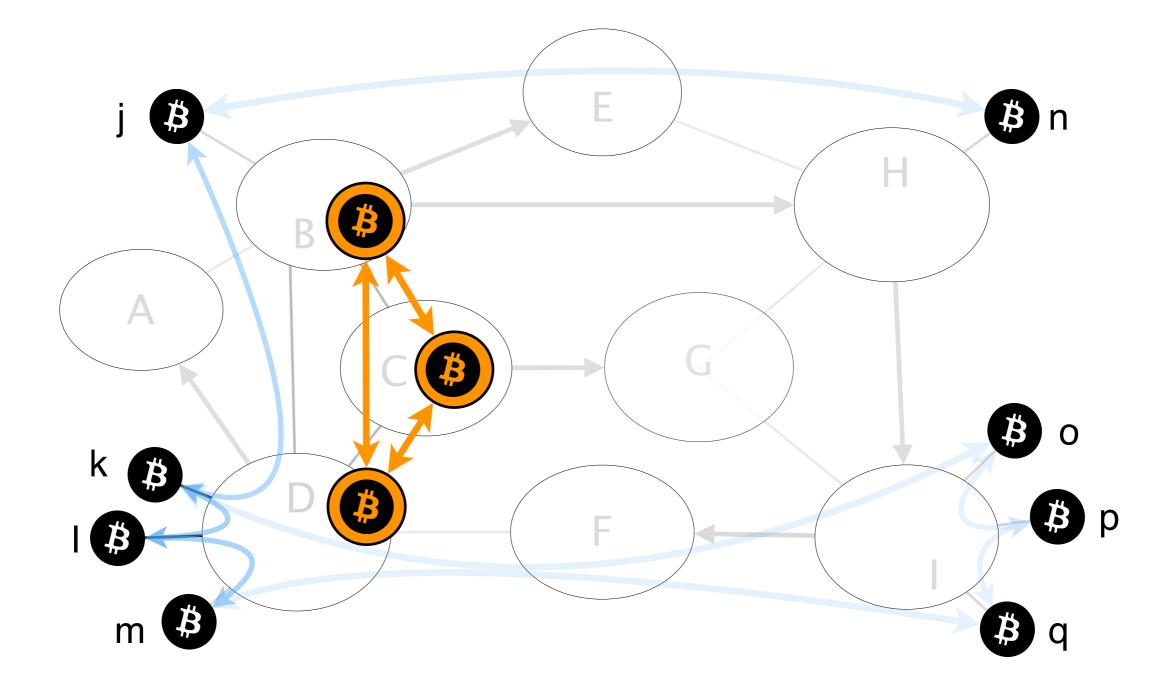
SABRE does not affect any of the regular Bitcoin clients



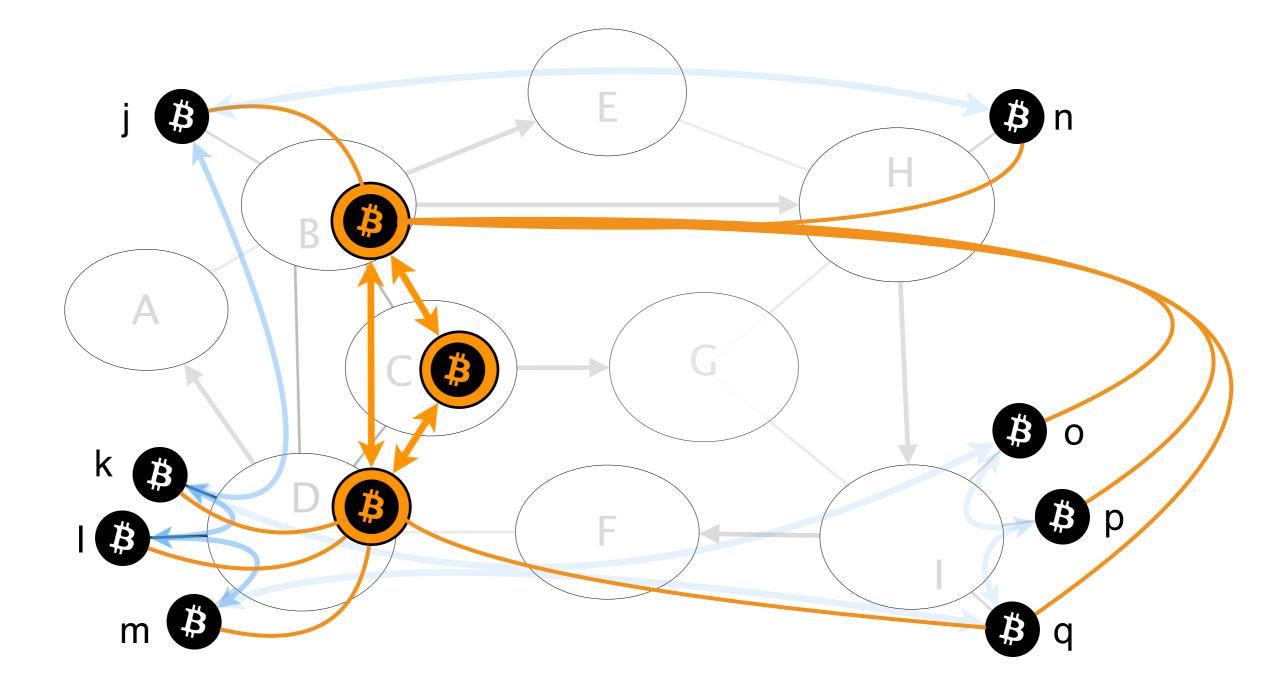
SABRE is an overlay network of special Bitcoin clients



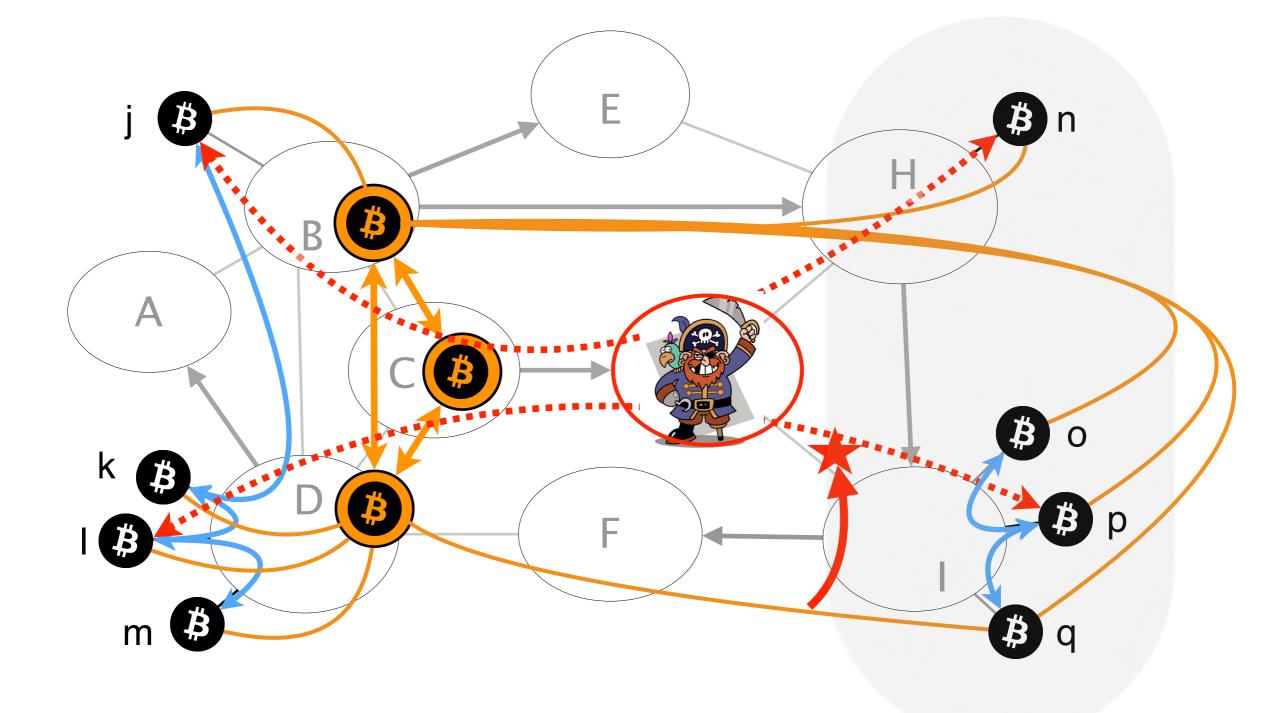
SABRE nodes are connected to each other



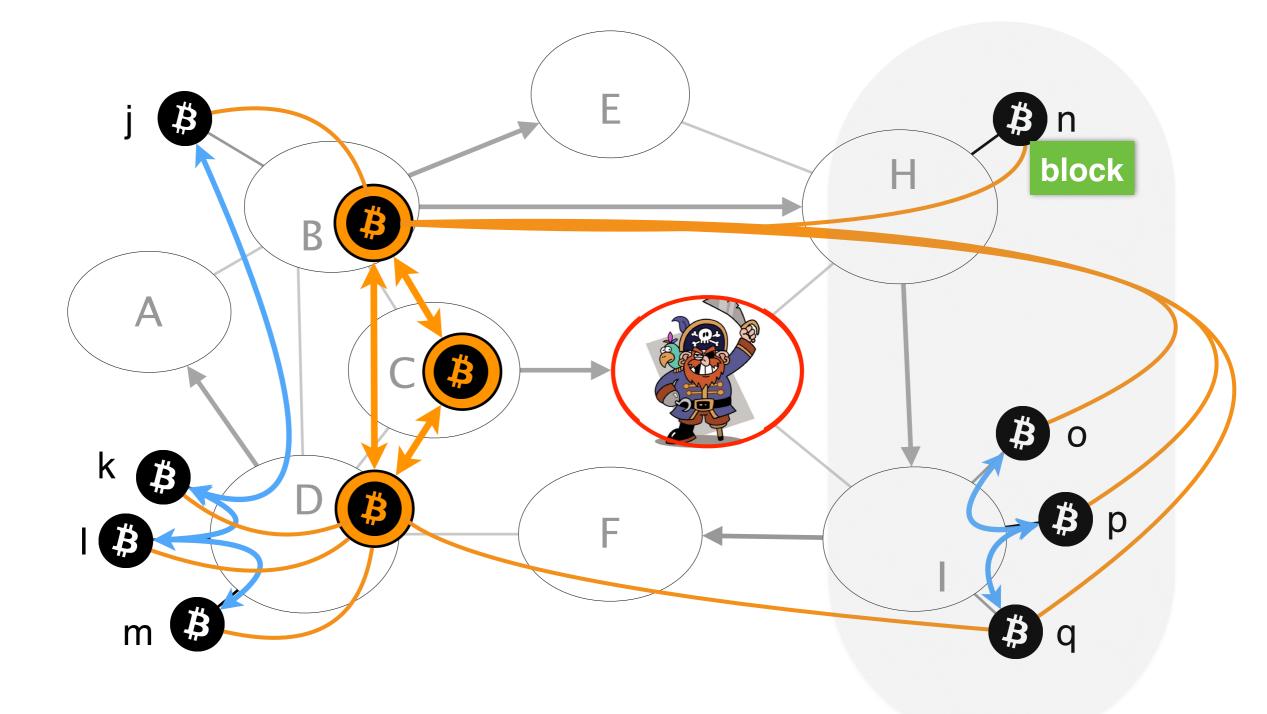
Each Bitcoin client connects to at least one SABRE node



SABRE protects the Bitcoin network from partition attacks



Block is propagated via the SABRE network



The attacker might try to fight back by attacking SABRE itself



The attacker might try to fight back by attacking SABRE itself

Attacker knows SABRE's locations and code

BGP hijacks against SABRE nodes



malicious requests to take down SABRE nodes



SABRE needs to...

secure relay-to-relay connections



SABRE needs to...

- **secure** relay-to-relay connections
- remain reachable by Bitcoin clients



SABRE needs to...

- secure relay-to-relay connections
- remain reachable by Bitcoin clients
- relay blocks under any load



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Network Design

Node Design



SABRE Protecting Bitcoin against Routing Attacks



SABRE location inherently safe locations

SABRE design software/hardware

Deployability deployment opportunities

SABRE Protecting Bitcoin against Routing Attacks



SABRE location inherently safe locations

SABRE design software/hardware

Deployability deployment opportunities

SABRE needs to...

- secure relay-to-relay connections
- remain reachable by Bitcoin clients

relay blocks





SABRE needs to...

secure relay-to-relay connections

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relay blocks under any load





SABRE selects nodes that satisfy three properties

each node is hosted in /24 IP prefixes

nodes are connected via financially & distance-wise optimal paths

relay graph is k-connected

SABRE selects nodes that satisfy three properties

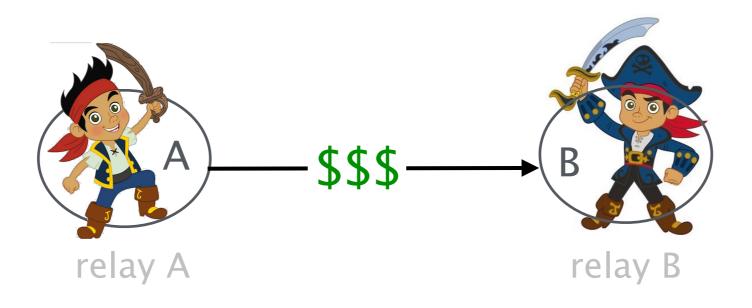
each node is hosted in /24 IP prefixes <

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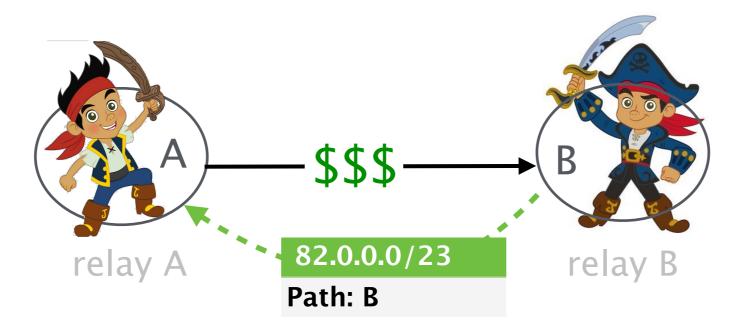
relay graph is k-connected

longer prefix hijacks are not possible

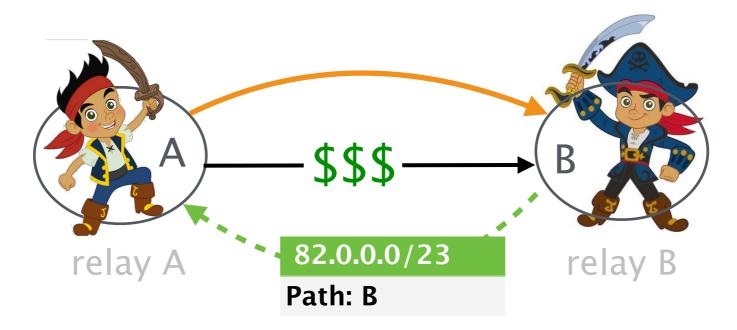
Relays A and relay B are hosted in ASes with customer-provider relationship



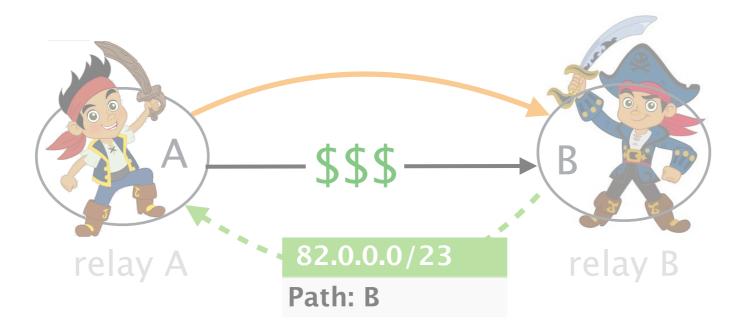
AS A receives a BGP advertisement from AS B for the prefix of relay B



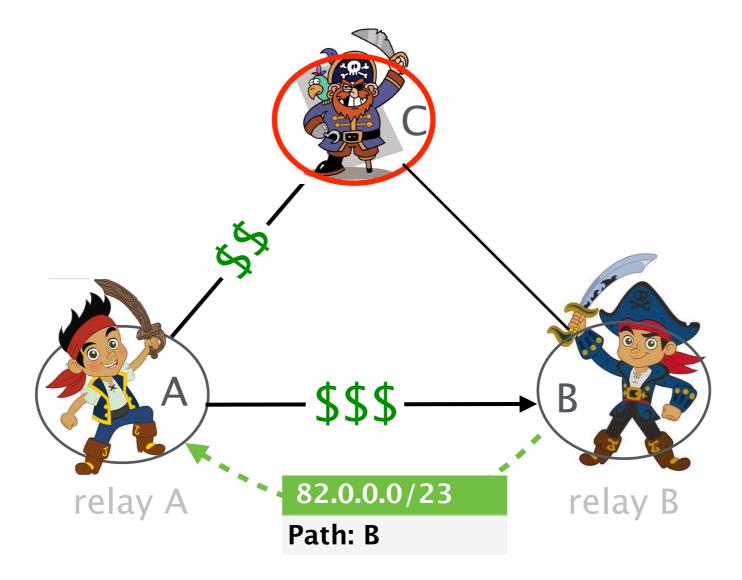
Relay A sends to relay B via a direct expensive link



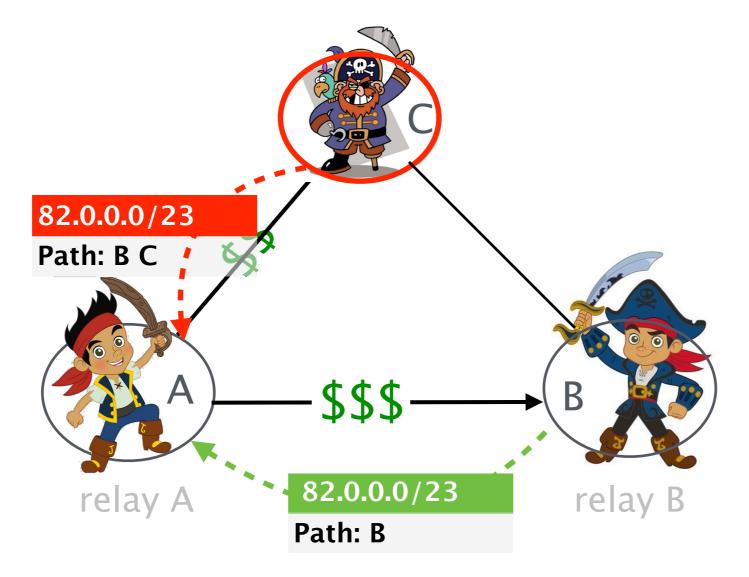
BGP is a policy-based protocol, with cost playing an important role



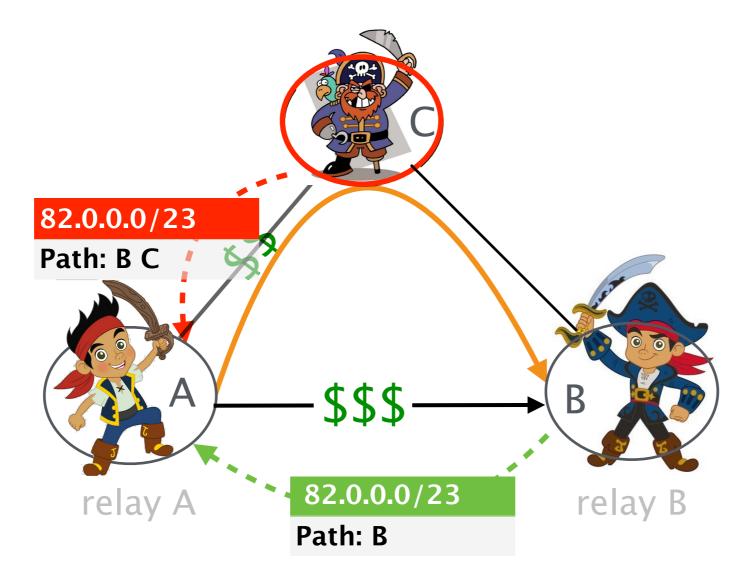
AS A has a malicious or compromised neighbor AS with a least expensive link



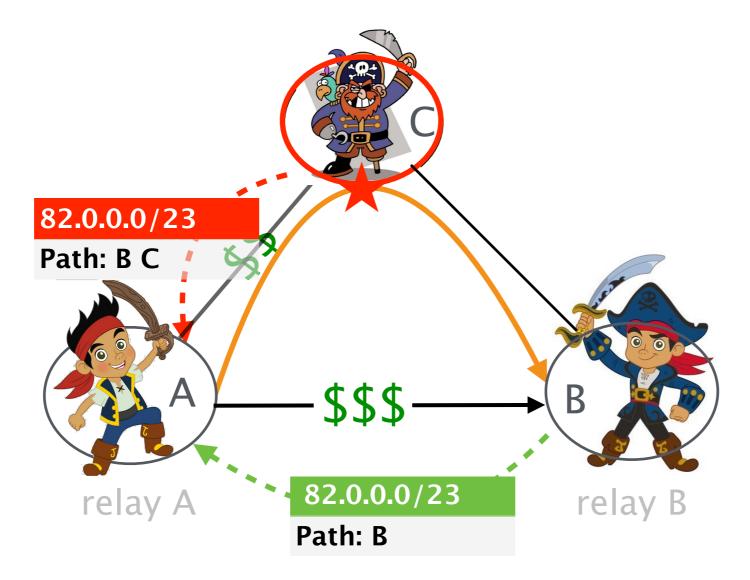
Attacker advertises AS B's prefix to AS A



AS A prefers the path via the attacker, because it is less expensive



The attacker can **disconnect** the relays



SABRE selects nodes that satisfy three properties

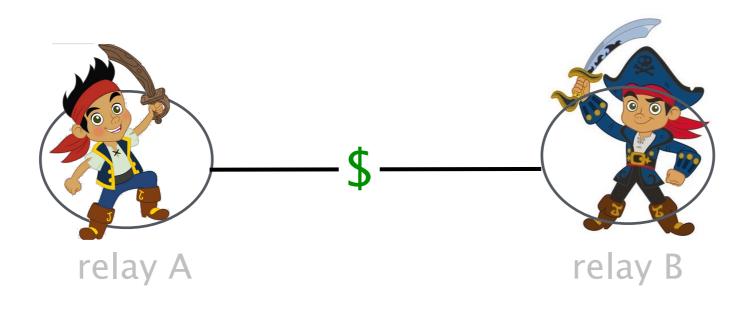
each node is hosted in /24 IP prefixes

nodes are connected via financially & distance-wise optimal paths

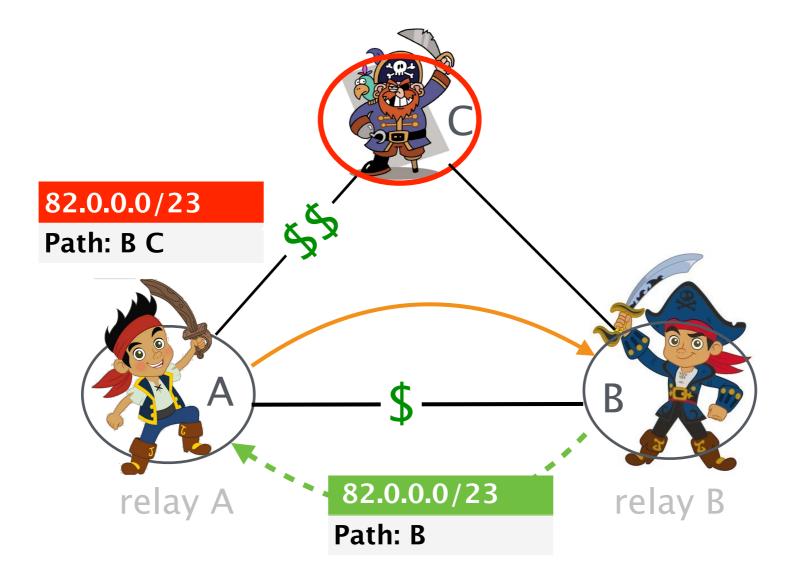
no strictly more preferred path exists

relay graph is k-connected

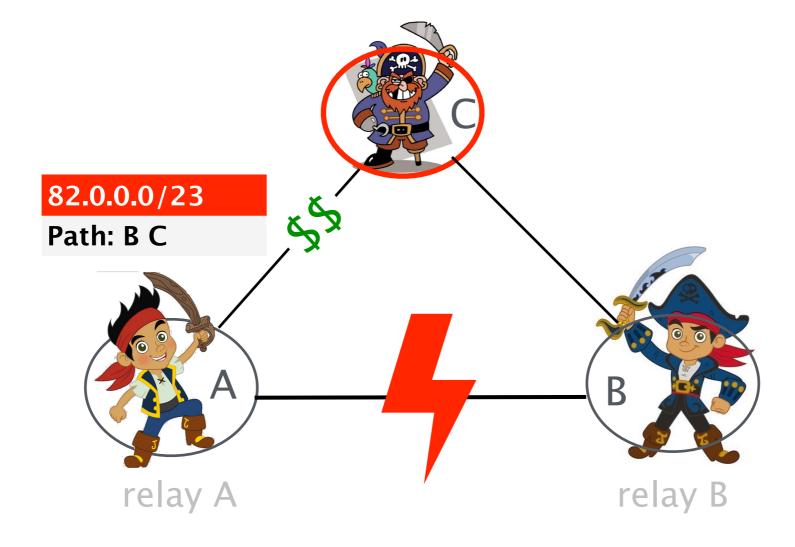
Relays A, B are hosted in ASes with a more cost effective agreement



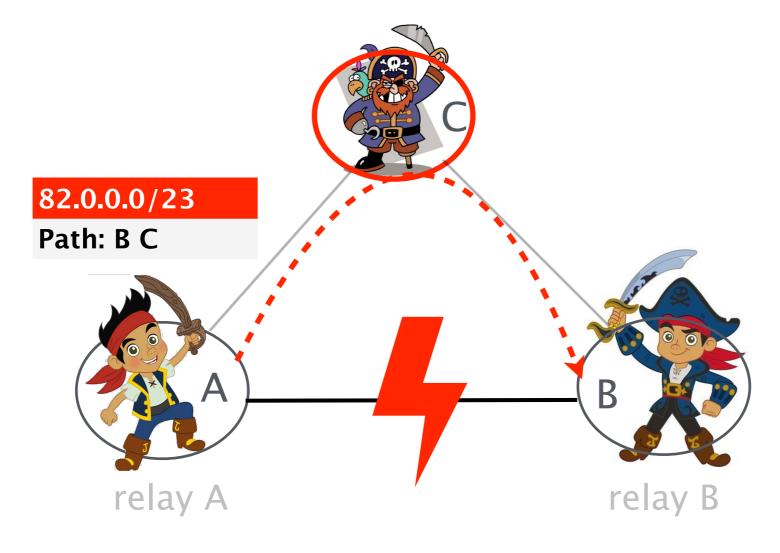
Attacker's advertisement is less preferred, thus attacker cannot discontent the relays



Aggreements can be revoked, link can be cut ...



Peering agreement can be revoked, link can be cut ... Relay A will inevitably send traffic via ASC



SABRE selects nodes that satisfy three properties

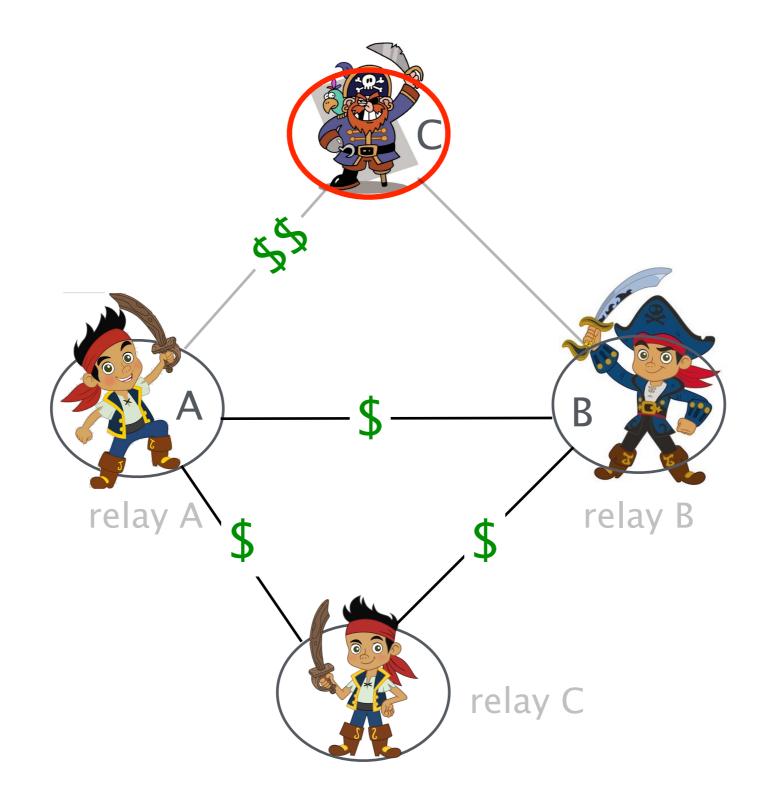
each node is hosted in /24 IP prefixes

nodes are connected via financially & distance-wise optimal paths

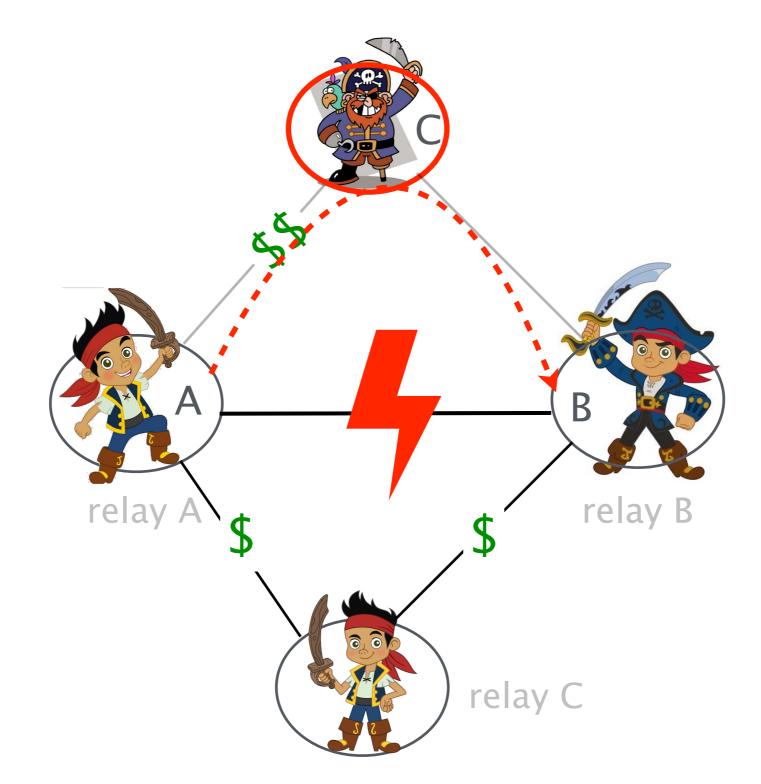
relay graph is k-connected

relay connectivity is not disrupted by any k-1 cuts

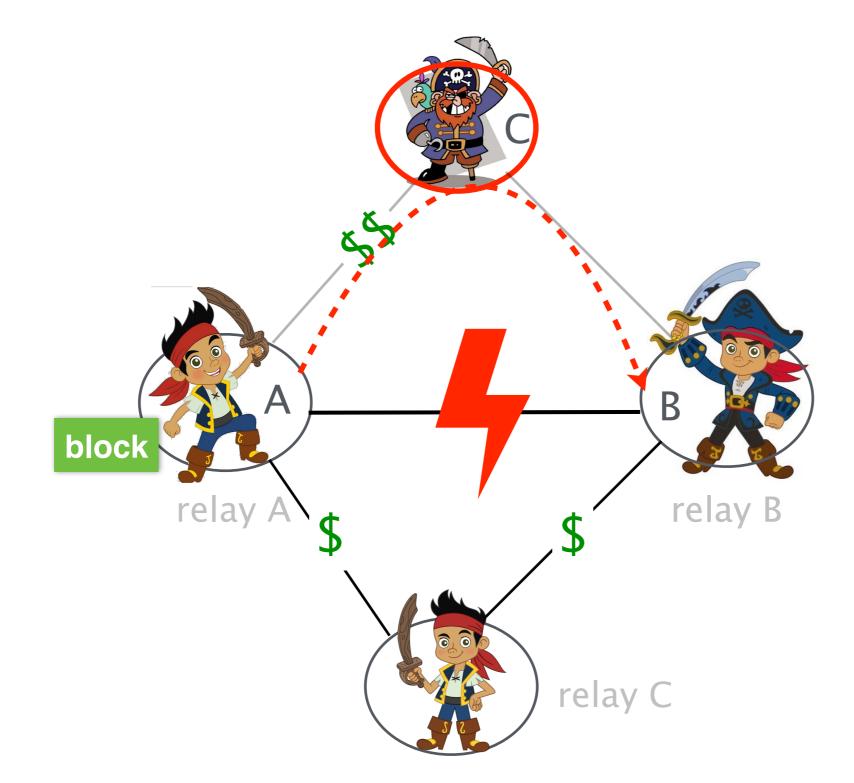
2-k connected graph retains connectivity even if one peering link is cut



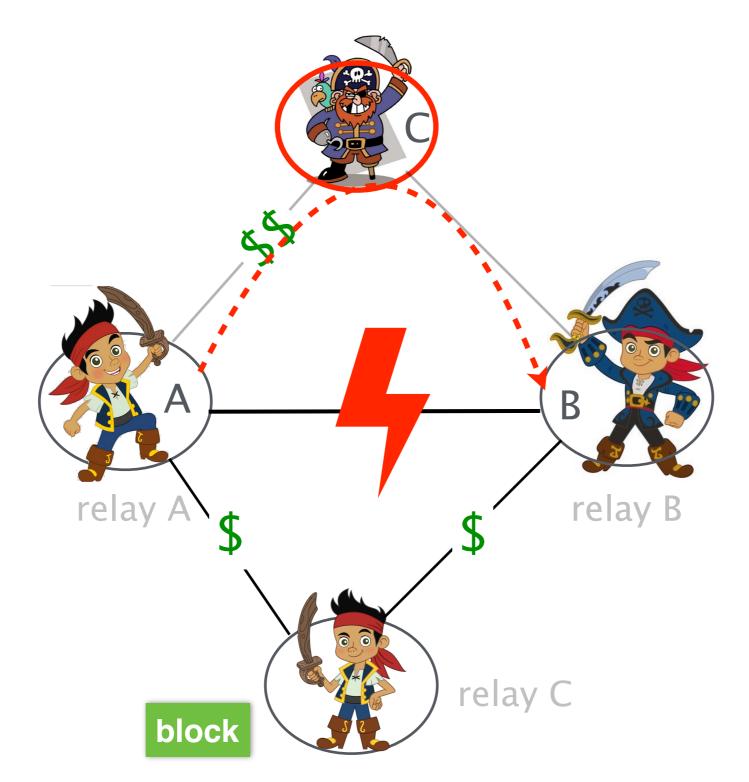
If the link between relays A and B is cut



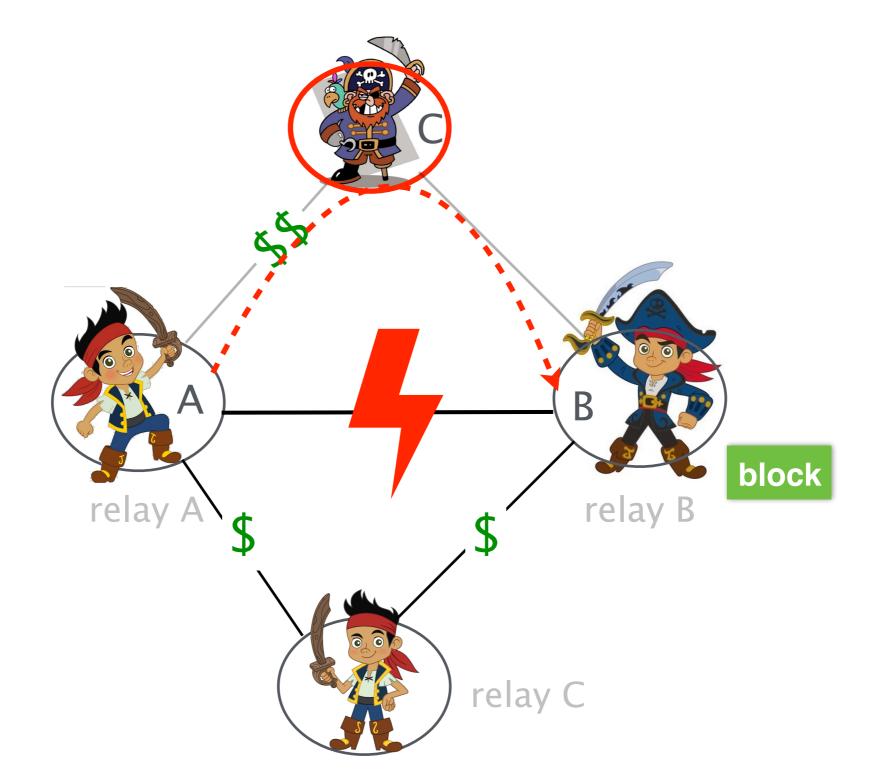
If the link between relays A and B is cut Relays A, B can still exchange blocks via the relay C



If the link between relays A and B is cut Relays A, B can still exchange blocks via the relay C



If the link between relays A and B is cut Relays A, B can still exchange blocks via the relay C



SABRE is an additional overlay network which allows communication, even if the Bitcoin network is partitioned

SABRE needs to...

- secure relay-to-relay connections
- remain reachable by Bitcoin clients
- relay blocks





SABRE positions nodes s.t. most clients are protected from each potential attacker by at least one relay node

see paper for more details

SABRE is an additional overlay network which allows communication, even if the Bitcoin network is partitioned

SABRE needs to...

- secure relay-to-relay connections
- remain reachable by Bitcoin clients
- relay blocks under any load

Network Design

Node Design



Network-wide attacks Node-level attacks

Network-wide attacks Node-level attacks

Network-wide attacks Node-level attacks

Network-wide attacks

What is the largest partition each single AS can create?

Node-level

attacks

How many clients are protected against isolation?

What is the largest partition each single AS can create?

What is the largest partition each single AS can create?

current network

any single AS in the world can create partitions of 90% of the clients

What is the largest partition each single AS can create?

current network

any single AS in the world can create partitions of 90% of the clients

6 SABRE nodes
 3-connected

only 3% of ASes in the world can create partitions of 15% of the clients

see paper for more results

Network-wide attacks

What is the largest partition each single AS can create?

Node-level attacks

How many clients are protected against isolation?

How many clients are protected against isolation?

How many clients are protected against isolation?

current network

at most 10% of Bitcoin clients are protected from 50% of ASes How many clients are protected against isolation?

current network

at most 10% of Bitcoin clients are protected from 50% of ASes

6 SABRE nodes
 5-connected

89.5% of Bitcoin clients are protected from 92.5% of ASes

see paper for more results

SABRE Protecting Bitcoin against Routing Attacks



SABRE location inherently safe locations

SABRE design software/hardware

Deployability deployment opportunities SABRE is an additional overlay network which allows communication, even if the Bitcoin network is partitioned

SABRE needs to...

- secure relay-to-relay connections
- remain reachable by Bitcoin clients
- relay blocks under any load



Two ways to deploy a SABRE node

Private deployment

Public deployment

Two ways to deploy a SABRE node

Private deployment

Public deployment

Serving few predefined clients

Private SABRE nodes need not scale

SABRE nodes need to

- establish connection to a predefined set of IPs
- receive and relay blocks
- be unreachable for unknown clients

Private SABRE nodes need not scale

SABRE nodes need to

- establish connection to a predefined set of IPs
- receive and relay blocks
- be unreachable for unknown clients

current Bitcoin client implementation hosted in a VM is sufficient

Two ways to deploy a SABRE node

Private deployment

Serving few predefined clients

Public deployment

Serving all Bitcoin clients

Public SABRE nodes need to scale

SABRE nodes need to

- maintain thousands of connections
- receive, verify and relay blocks fast
- protect against spoofing and malicious request

Public SABRE nodes need to scale

SABRE nodes need to

- maintain thousands of connections
- receive, verify and relay blocks fast
- protect against spoofing and malicious request

Simple software implementation would not suffice

SABRE can leverage programmable data planes SABRE DP

is faster than any server optimization

can serve few Billions of packets per second

NetChain: Scale-Free Sub-RTT Coordination NDSI 2018

is faster than any server optimization

protects against malicious requests

Dynamic Black/White lists anti-spoofing mechanism & DoS protection

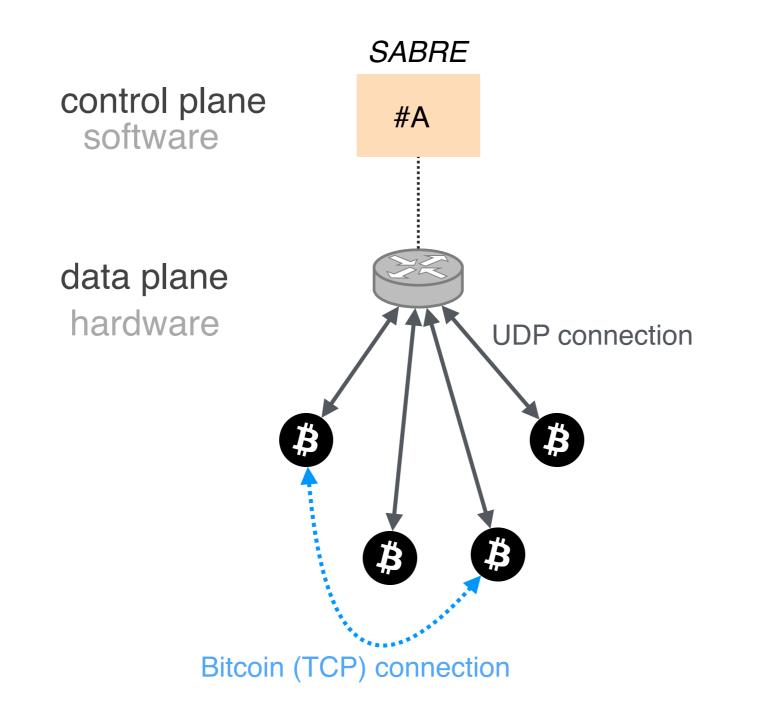
is faster than any server optimization

protects against malicious requests

minimum software interaction

almost all clients served directly from hardware Not all operations can be done in hardware

Not all operations can be done in hardware SABRE node has both software and hardware parts



SABRE Protecting Bitcoin against Routing Attacks



SABRE location inherently safe locations

SABRE design software/hardware

Deployability deployment opportunities

bootstrap with a software-only SABRE

decreased cost allows private deployments

bootstrap with a software-only SABRE

multiple SABRE relays can co-exist

each party (e.g. pool) can deploy their own SABRE without coordination

bootstrap with a software-only SABRE

multiple SABRE relays can co-exist

community's consensus is not required

clients can connect to both relays and regular clients

bootstrap with a software-only SABRE

multiple SABRE relays can co-exist

community's consensus is not required

network design applies to other relays

e.g., FIBRE, FALCON can relocate relays following SABRE location algorithm

SABRE Protecting Bitcoin against Routing Attacks



SABRE location inherently safe locations

SABRE design software/hardware

Deployability deployment opportunities

SABRE Protecting Bitcoin against Routing Attacks



SABRE can protect Bitcoin from partitions by placing few relay nodes in selected locations

SABRE can operate seamlessly under high load by serving clients directly in hardware

SABRE can be partially deployed and benefit early adopters e.g., each pool can deploy SABRE in software